

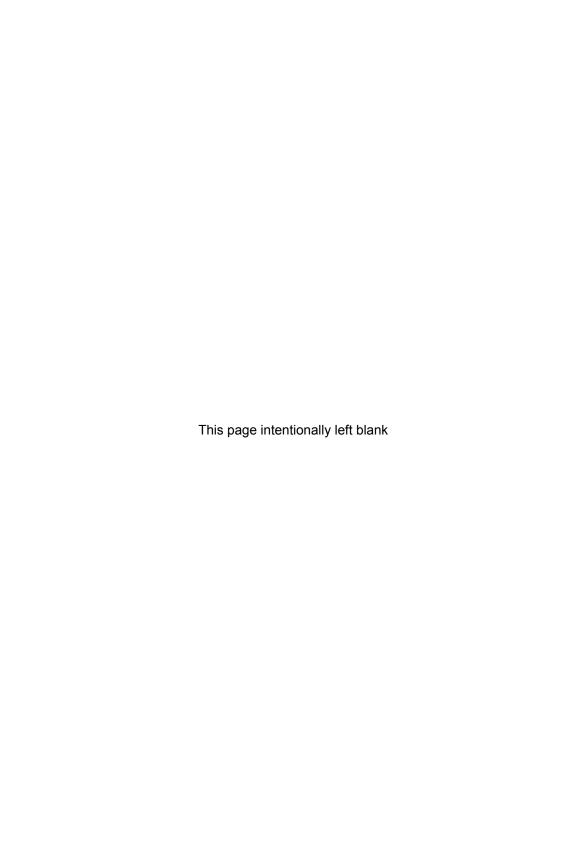
Petri Nets

Fundamental Models, Verification and Applications

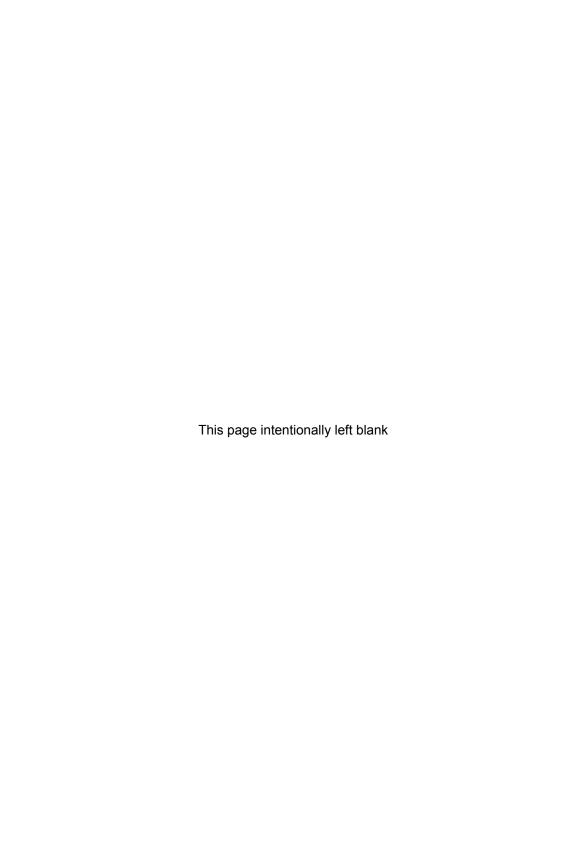
Edited by Michel Diaz











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Table of Contents

Preface	χv
Introduction	vii
PART 1. FUNDAMENTAL MODELS	1
Chapter 1. Basic Semantics	3
1.1. Automata or state machines	3
1.1.1. Automata and state machine models	3
1.1.2. Tasks and processes	5
1.1.3. Some models	6
1.2. State machines and Petri nets (PN)	8
1.2.1. Composing state machines	8
1.2.2. Composition and synchronization	10
1.3. Concepts and definitions	11
1.3.1. Local states and enabling	12
1.3.2. Definition of the semantics of parallelism	12
1.3.3. Firing transitions	13
1.4. Accessibility graph or marking graph	15
1.5. Some basic models	18
1.5.1. Co-begin (parallel start) and co-end (synchronized termination) .	18
1.5.2. Synchronization by a signal	18
1.5.3. Mutual exclusion	19
1.5.4. The reader and writer mechanisms.	21
	23
1.5.5. Bounded buffers	
1.6. Conclusion	24
1.7. Bibliography	2.5

Chapter 2. Application of Petri Nets to Communication Protocols 27 Michel DIAZ
2.1. Basic models
2.2. A simple establishment of a connection
2.2.1. Different global semantics
2.2.2. Conclusion
2.3. The alternating bit protocol (ABP): model and verification
2.3.1. Loss of messages
2.3.2. Modeling losses
ϵ
2.5. Conclusion 37 2.6. Bibliography 38
2.0. Bioliography
Chapter 3. Analysis Methods for Petri
Serge HADDAD and François VERNADAT
3.1. Introduction
3.2. Behavioral analysis of Petri nets
3.2.1. Semantics of a net
3.2.2. Usual properties
3.3. Analysis of nets by linear invariants
3.3.1. Definitions and first applications
3.3.2. Flow computations
3.3.3. Semiflow computation
3.3.4. Application of invariants to the analysis of a net
3.4. Net reductions
3.4.2. Post-agglomeration of transitions
3.4.3. Deletion of redundant places
3.5.The graph of a Petri net
3.5.1. General results
3.5.2. State machines
3.5.3. Event graph
3.5.4. Free choice net
3.6. Bibliography
Chapter 4. Decidability and Complexity of Petri Net Problems 87 Serge HADDAD
4.1. Introduction
4.2. Decidability and complexity notions
4.3. Theoretical results about the reachability graph
4.4. Analysis of unbounded Petri nets

4.4.1. Construction of the covering graph	95
4.4.2. Shortest sequences	101
4.4.3. Backward analysis	103
4.5. The reachability problem	105
4.5.1. A necessary condition for reachability	106
4.5.2. A sufficient condition for reachability	106
4.6. Extensions of Petri nets	109
4.6.1. Netswith inhibitor arcs	109
4.6.2. Self-modifying nets	111
4.6.3. Recursive nets	113
4.7. Languages of Petri nets	116
4.8. Bibliography	120
Chapter 5. Time Petri Nets	123
Bernard Berthomieu, Marc Boyer and Michel DIAZ	
5.1. Introduction	123
5.2. Time Petri nets	126
5.2.1. Time nets	126
5.2.2. States and firing rule	127
5.2.3. Set of states, schedules	128
5.2.4. Firing domains	129
5.3. Behavior characterization – state classes' method	130
5.3.1. State classes	130
5.3.2. Transitions between state classes	131
5.3.3. State class equality	134
5.3.4. Class graph	136
5.3.5. Marking graph and class graph	137
5.4. Analysis – operating the state class graph	138
5.4.1. Analyzing behavior of time-dependent systems	138
5.4.2. Marking reachability	139
5.4.3. Boundedness	139
5.4.4. Specific properties for set of markings or firing sequences	143
5.4.5. Time-dependent analyses, existence of schedules	143
5.5. Application example	144
5.6. Extensions and variations	149
5.6.1. Interpreting multi-enabled transitions	149
5.6.2. Other time extensions	154
5.7. Implementation using the Tina tool	156
5.7.1. Tina tool	156
5.7.2. Application example	156
5.8. Conclusion	158
5.9. Bibliography	159

Chapter 6. Temporal Composition and Time Stream Petri Nets
6.1. Time, synchronization and autonomous behaviors
6.2. Limitation of time PNs
6.3. Temporal composition
6.4. Temporal composition and temporal synchronization
6.4.1. The semantics of "waiting"
6.4.2. Pragmatics and time assumptions
6.5. Time stream PNs
6.5.1. Definition of the model
6.5.2. The different firing semantics
6.5.3. Relating times behavior
6.5.4. TSPN with structured streams
6.6. Application to multimedia systems
6.6.1. Jitter in streams
6.6.2. Intra- and inter-stream drifts
6.6.3. Modeling stream composition
6.6.4. Principle of modeling multimedia systems
6.6.5. Modeling multimedia scenarios
6.6.6. TSPN for designing hypermedia architectures
6.7. Conclusion
6.8. Bibliography
Chapter 7. High Level Petri Nets
Claude GIRAULT and Jean-François PRADAT-PEYRE
7.1. Introduction
7.2. Informal introduction to high level nets
7.2.1. A client-server model
7.2.2. Client distinction
7.2.3. Server distinction
7.2.4. Equivalent unfolded net
7.2.5. Colored model for the alternate bit protocol
7.3. Colored net definition
7.3.1. Notation
7.3.2. The formalismof colored nets
7.3.3. Unfolding of a colored net
7.4. Well-formed net definition
7.4.1. Color domains
7.4.2. Color functions
7.4.3. Guards
7.4.4. The formalismofwell-formed nets
7.4.5. Regular nets and ordered nets

7.5. Other high level formalisms 7.5.1. Interpreted nets 7.5.2. Algebraic nets 7.6. Conclusion 7.7. Bibliography	212 212 214 219 219
Chapter 8. Analysis of High Level Petri Nets	221
8.1. Introduction 8.2. The symbolic reachability graph 8.2.1. Symbolic markings 8.2.2. Symbolic marking representation 8.2.3. Symbolic firing rule 8.2.4. Example of a symbolic reachability graph 8.2.5. Properties of the SRG 8.3. Colored invariants 8.3.1. Definition of invariants of high level Petri nets 8.3.2. Computing flows of a high level net: principles and difficulties 8.3.3. Computing a non-parametrized generative flow family 8.3.4. Parametrized generative family of flows for regular nets 8.3.5. Computation of positive flows 8.4. Structural reductions 8.4.1. Principles of extension to high level nets 8.4.2. Pre-agglomeration and post-agglomeration of transitions 8.4.3. Deletion of an implicit place 8.4.4. Application examples 8.5. Conclusion 8.6. Bibliography	221 222 223 225 230 2344 238 242 243 250 252 253 254 255 261 265 266
Chapter 9. Stochastic Petri Nets	269
Serge HADDAD and Patrice MOREAUX	
9.1. Introduction 9.2. A stochastic semantics for discrete event systems 9.2.1. The stochastic model 9.2.2. Analysis with renewing theory 9.2.3. Discrete time Markov chains 9.2.4. Continuous time Markov chains 9.2.5. Semi-Markovian processes 9.2.6. Regenerative Markovian processes 9.3. Stochastic Petri nets 9.3.1. Stochastic Petri nets with general distributions	269 271 271 273 274 276 278 279 280 280
9.3.2. Stochastic Petri nets with exponential distributions	282

x Petri Nets

11.4.4. A tensor decomposition method for SWN 11.4.5. Application in the asynchronous case 11.5. Conclusion 11.6. Bibliography	338 340 344 344
PART 2. VERIFICATION AND APPLICATION OF PETRI NETS	347
Chapter 12. Verification of Specific Properties	349
12.3. Temporal logic	349 352 354 354 357 369 371 375 385 395 401 401 407 411
Chapter 13. Petri Net Unfoldings – Properties	415
13.2. Elementary concepts. 13.2.1. Preliminary information 13.2.2. Net homomorphisms 13.2.3. Occurrence nets 13.3. Branching processes and unfoldings 13.3.1. Branching processes 13.3.2. Unfoldings 13.4. Finite prefixes. 13.4.1. Definition. 13.4.2. Adequate orders and complete finite prefixes 13.4.3. Verification of safety properties 13.4.4. Detection of infinite behaviors	415 416 416 417 418 421 423 424 424 425 427 428
	432 432

Chapter 14. Symmetry and Temporal Logic	435
Serge HADDAD and Jean-Michel ILIÉ 14.1. Introduction	435 437 437 439 441 444 452 454 457 458 459
Chapter 15. Hierarchical Time Stream Petri Nets	461
15.1. Introduction . 15.2. Structured time stream Petri nets 15.2.1. Motivations. 15.2.2. From composition to abstraction . 15.2.3. Verification of temporal coherence 15.3. Combining abstraction and temporal composition . 15.3.1. Modularity and abstraction of temporal behaviors . 15.3.2. Hierarchical time stream Petri nets . 15.3.3. Interrupts as a combination of abstraction and temporal composition . 15.3.4. HTSPN state . 15.4. Examples . 15.4.1. Modeling hypermedia systems . 15.4.2. A solution to "lip-synchronization" using HTSPNs . 15.5. Conclusion . 15.6. Bibliography .	461 462 462 463 466 468 468 471 472 474 474 475 477 478
Chapter 16. Petri Nets and Linear Logic	481
16.1. Introduction	481 483 483 483 484

16.3. Petri nets and linear logic	485
16.3.1. Various approaches	485
16.3.2. Approach with marking	486
16.3.3. Equivalence between reachability in a Petri net and	
provability of one sequent in linear logic	488
16.4. Sequent labeling and graph of precedence relations	490
16.4.1. Labeling	490
16.4.2. Graph of precedence relations	491
16.4.3. Conflicts of transitions and tokens	492
16.5. Temporal evaluation of scenarios	493
16.5.1. Introduction	493
16.5.2. Simple temporal networks	494
16.5.3. Temporal labeling	496
16.6. Conclusion	499
16.7. Bibliography	500
Chapter 17. Modeling of Multimedia Architectures: the Case of	
Videoconferencing with Guaranteed Quality of Service	501
Philippe OWEZARSKI and Marc BOYER	
17.1. Introduction	501
17.1. Introduction:	502
17.2.1 Multimedia information: characteristics and requirements	502
17.2.2. Asynchronous distributed systems	505
17.2.2. Asynchronious distributed systems	507
17.3.1. Modeling requirements	507
17.3.2. Modeling example for a videoconference application	510
17.3.2. Modeling example for a videoconference application	512
17.4.1. Introduction	512
17.4.2. Modeling of a videoconference application.	512
17.4.2. Modeling of a videoconference application	516
17.5. Conclusion	522
	523
17.6. Bibliography	323
Chapter 18 Performance Evaluation in Manufacturing Systems	527
Isabel Demongodin, Nathalie Sauer and Laurent Truffet	321
18.1. Introduction	527
18.2. Modeling of manufacturing systems	528
18.2.1. Discrete event systems aspects	529
18.2.2. Cyclic aspects	533
18.2.3. High throughput aspects	538
18.2.4. Weighted marked graphs	543
18.3 Evaluation of manufacturing systems	544

xiv Petri Nets

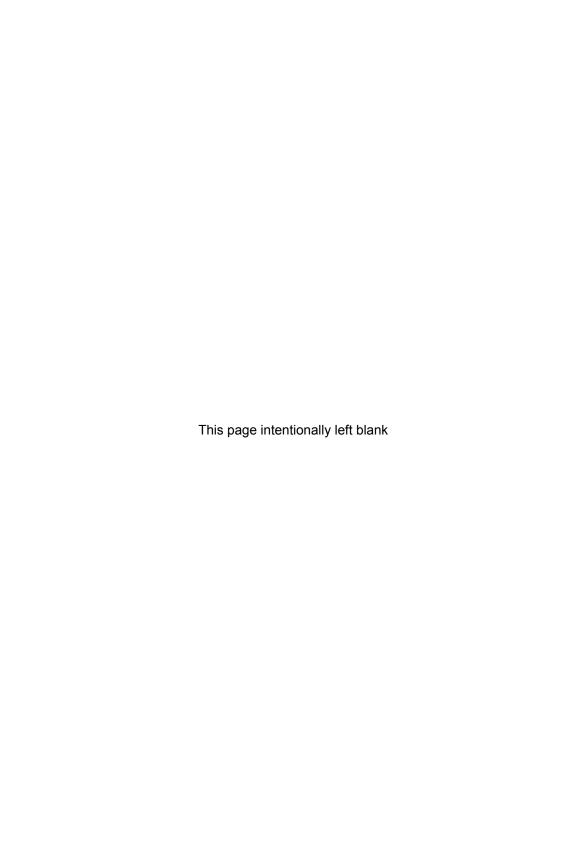
18.3.1. Performance evaluation methods	544
18.3.2. Deterministic and stochastic discrete marked graphs	549
18.3.3. Discrete weighted marked graphs	551
18.3.4. Continuous weighted marked graphs	552
18.4. Optimization of manufacturing systems	559
18.4.1. Deterministic marked graphs	560
18.4.2 Stochastic marked graphs	561
18.4.3. Extension to deterministic weighted marked graphs	563
18.4.4. Applications	567
18.5. Conclusions	571
18.6. Bibliography	572
Conclusion	577
List of Authors	579
Index	581

Preface

Future advanced architectures, such as embedded systems, having a greater complexity and new quality requirements, will need a more precise specification and better control of their design process. In order to acquire the corresponding fundamental knowledge, it is essential to rely upon approaches based on the use of adequate system models. In particular, such approaches need to acquire a deep understanding of the system, including its local behaviors and its communications, based on a well-defined representation of the designed architecture. This representation should be used as early as possible to analyze and validate the design. The goal of this volume is to present a family of formal specification models, based on Petri nets and extensions of Petri nets, because they are defined by simple and clear semantics, allow easy modeling of system key mechanisms, and are supported by strong analysis methods and tools. Furthermore, this set of models can be used for all design aspects, i.e. to specify functional behaviors, and to include temporal or stochastic requirements.

The main results related to this approach are given in this volume, in two parts, one presenting the fundamental models, and the other being dedicated to verification and applications. We have tried to highlight the important characteristics and the main properties of these models, and to show how they lead to the emergence of a full design methodology, which is both complete, in terms of all possible functional and other analysis, and integrated, because the same basic semantics are used for the full design support. We think that this volume should greatly help any designer to build the new forthcoming generation of distributed systems.

Lastly, I would like to thank all the authors who contributed to this book, for their expertise, their seriousness, their technical inputs, and for the great job they have done.



Introduction

New technologies in processors and networks allow system designers to conceive and build advanced and sophisticated parallel and distributed architectures, which need to integrate non-functional real time and stochastic constraints with functional distributed processing and communication.

The global behavior of such systems depends first on the local activities and data, but also on the messages sent and received by the various interconnected subsystems. As a matter of fact, understanding, expressing, specifying and validating such global behaviors proves to be a problem of very high complexity, leading to many design and implementation difficulties and bugs. For example, when considering n connected processors, they can run, at a given instant in time, using 2×2 , 3×3 communications, etc., or a full communication, in which all n processors interact. The sum of the resulting combinations, of the order of 2^n , shows the complexity of the resulting conceptual problems, and explains in particular the increasing difficulty obtained when passing from an interconnection of a few processors to an interconnection of a large number of processors: when the number of processors varies from 2 to 10, the difficulty coefficient goes from 4 to about 1,000.

It should then be clearly understood that designing such distributed architectures leads to a very complex conceptual task, which has to be based on a well-defined methodology to be able to manage all system requirements and behaviors.

Design and specification

The design process starts by giving the different functions and agents which are required, and the way they are structured; second, the designers define the behaviors of the various processes and entities, and the way they communicate; then, if they want to analyze the correctness of the design as soon as possible, an adequate

approach is needed to represent, in an explicit way, the (full) system global behavior, in particular to be able to check potential unanticipated sub-behaviors.

To check the design correctness, it is essential to use a precise *model* of all critical mechanisms, functions, sub-systems, etc., and then, whenever possible, to use a *formal model*, to define a mathematical representation of the system. Checking the *correctness* validation of the design at this step is then conducted by checking the behavior of this system model.

Note that, after a given adequate sequence of more or less formal validation steps based on models, the system will be defined as 'fully designed' and will be implemented using adequate tools and languages.

Formal approaches have been used for many years for the verification of communication protocols. Two principal approaches have been used., i.e. basic formal models, such as automata, Petri nets, process algebras, etc., and formal description techniques for protocols, such as Estelle, LOTOS, SDL, etc.

This volume proposes and develops a design and validation methodology that relies on the use of a family of basic formal models that are rather easy to understand, and able to:

- describe the semantics of all basic building mechanisms;
- clearly specify the interconnection and communication semantics;
- unambiguously describe the resulting behaviors;
- validate the system during the first phases of its design by using support tools.

In general, basic non-language oriented graphical models, that do not include language-specific operators and statements, lead to the simplest solutions for representing basic mechanisms in a very abstract and integrated way.

For this reason too, this volume selected a basic, language-independent set of models to represent and manipulate the fundamental concepts of communicating architectures.

Selecting a model

Several models exist, and each model has particular characteristics, more or less relevant for a specific design. Consequently, the choice of the right model depends on the designed system and on the properties to be analyzed, as the model must be able to describe the design, and also to allow the designer to check the validity of the required properties.

In general, the designer must have a good understanding of the fundamental semantics of the system, i.e. of its basic building mechanisms. Thus, for simple architectures, modeling will be able to represent in a faithful way all details of the system. However, for complex systems, it will generally be impossible, for economic reasons, to represent the details of all existing functions, and it will become necessary to select and validate certain building blocks, i.e those most likely to lead to erroneous behaviors.

Of all existing models, *Petri nets* (PN) and their extensions are of undeniable fundamental interest, because they:

- provided the first modeling approaches for the semantics of concurrent systems, and were used to model the behaviors of the first parallel and distributed basic mechanisms:
- define easy graphic support for the representation and the understanding of these basic mechanisms and behaviors:
- prove to be, starting from state machines, an easy extension of previous approaches and handle, at the same time, the creation and the analysis of models;
- express very simply the main basic concepts in communication, including waiting and synchronization, and furthermore take into account their temporal and stochastic parameters;
- ensure, being unrelated to a particular implementation language, the independence of the specification with respect to its implementation.

Furthermore, many validation methods have been developed, using a great number of theoretical results and support tools, able to manipulate functional, temporal, and stochastic behaviors. Finally, models based on PNs will help us to understand, define and analyze the behavior of these systems, in the preliminary and first steps of their design.

For all these reasons, a set of Petri net models was selected in this book to represent and manipulate the fundamental concepts of communicating architectures.

Petri nets

PNs were introduced by A.C. Petri in 1962 to synchronize communicating automata, and were then extended to define a large set of models, with increasing complexity and capabilities.

As will be seen, this family of PNs, starting from the simple traditional state machines, now allows system designers to handle in an integrated way the functional (qualitative) and the non-functional (e.g. quantitative) temporal and stochastic capabilities of systems.

Extensions of PNs were proposed according to two important axes:

- a) for *qualitative properties* and behaviors, to use simpler and more compact models, by high level PNs, for handling generic behaviors (e.g. individual) and data, predicates and functions;
- b) complementing this first axis, for *quantitative properties* and behaviors, to extend the previous models by integrating quantitative constructs and parameters related to temporal and stochastic requirements.

It is significant to note that all first and conceptual studies in these quantitative fields were carried out using PN-based models.

Functional qualitative properties

The first PN model, called the Condition-Event PN, was based on the use of Boolean values: true or false. It was generalized by Place-Transition PNs, now simply called PNs, which can use integers. This volume will begin with their presentation and validation.

Non-functional quantitative properties

The fundamental contributions of the second axis considers:

- time PNs, or TPNs, used for systems whose behaviors depend explicitly on temporal values;
- *stochastic PNs*, or SPNs, for which distributions are attached to the model, in particular for performance evaluation and reliability.

Families of PNs

When applied to the modeling of systems, it rapidly becomes apparent that these models do not have the same application power, in terms of:

- definition and description of the concepts for parallelism, distribution, and synchronization;
 - understanding and using the temporal and stochastic semantics;
- analyzing the possibly different mechanisms and behaviors, in very different contexts and applications.

Figure 1 represents some of the principal models of this family. In this figure, an arrow means that the model at the end of the arrow was proposed after the model at the beginning of the arrow, and so gives the steps followed by the research to propose and develop these principal PN-based models.

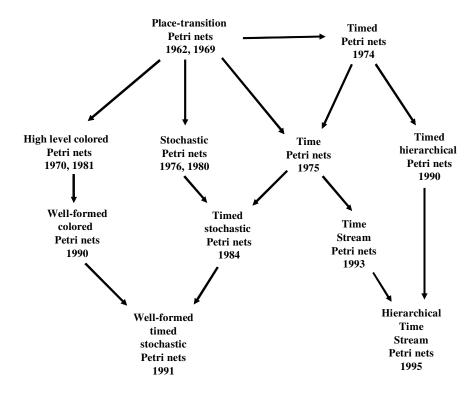


Figure 1. The main Petri net models

Figure 2 gives a more conceptual view of these models, by clarifying their syntactic and semantic relationships. In this figure, three fields are respectively defined by:

- a discrete state semantics, for non-temporal and non-stochastic nets, behaviors being represented by a finite graph of all model states;
- a semantics on continuous time, for extended behaviors based on dense time models:
 - and stochastic semantics, for behaviors including distributions.

Let us emphasize that these models have three models of reference, respectively PN, TPN, and SPN. Moreover, each model is a pure extension of a previous one, as it can by simplified to become a basic PN model.

As seen in the figure, the models derived from the reference models:

- lead to more compact models, i.e. are abbreviations, that do not increase the
 expressiveness of the model, but simplify the model and the system specification;
- or are more powerful in terms of expressive power, i.e. are able to describe mechanisms which could not be described by the unextended models (e.g. introducing time parameters, stochastic distributions, etc., for real-time or dependable systems).

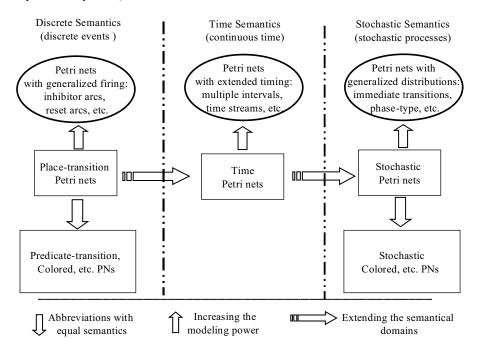


Figure 2. Semantics domains of the Petri net-based models

For example:

- PN led to PN with inhibitor arcs (to test the presence of zero token in one place), PN with reset arcs, etc.;
- TPN led to TPN with streams to compose and synchronize independent behaviors with independent temporal constraints, etc.;

- SPN led to SPN with immediate transitions in order to manage the case where transition cannot be delayed, etc.

Consequently, many different models exist, of different power and for different fields of application, but they follow the same semantics basis, and will allow the designers to carry out coherent complementary analyses to validate the correct operation of the modeled (and designed) systems.

The semantics of these models and their properties were used to select, define, and study the most important members of the PN family in the two parts of this volume.

Table of contents of the volume

Part 1 is dedicated to fundamental models and contains 11 chapters. Part 2 addresses verification and applications, and contains the last 7 chapters.

Part 1

Chapter 1 introduces Place-Transition PNs, more simply called Petri Nets (PNs). It gives their fundamental definitions, presents some basic models and clarifies their interest.

Chapter 2 illustrates an application in a very important area: communication protocols; simple PN examples show at the same time the power of the model and the interest of the formal analysis.

Chapter 3 first introduces the general properties that can be checked using PNs (blocking, reachability or accessibility, etc.) and the verification approach that uses the graph of the reachable (or accessible) states. The set of reachable (or accessible) states is the set of states that are reachable or accessible from a given initial state. Two optimization methods of analysis are then presented, one based on linear algebra techniques, and the other one exploiting the topological structure of the PNs.

Chapter 4 deals with the decidability and complexity problems related to checking these general properties.

Chapters 5 and 6 consider models and behaviors based on explicit values of time, and show how to model temporal mechanisms.

Chapter 5 presents the general model, Time PN or TPN, which associates a given interval (minimum, maximum) to each transition; this gives the first semantics for handling time and verifying temporal behaviors.

Chapter 6 presents a general model for composing temporal behaviors and systems. It gives the semantics of temporal composition by a new model, Time Stream PN, for composing autonomous (temporal) flows. It emphasizes their interests and applications for systems having independent temporal constraints, which sometimes interact.

Chapters 7 and 8 again consider PNs, i.e. non-temporal PN models, but define an abbreviation of a PN by a general model, which becomes able to represent, in a very compact way, a given set of similar parallel behaviors. The problems associated with this abbreviation are, on the one hand, to define a compact formalism and, on the other hand, to propose new validation techniques to handle this model, i.e. to avoid the obvious solution that consists of unfolding it into a very large PN.

Chapter 7 presents the main PN abbreviations, while concentrating on Colored PNs, which is the most frequently used model.

Chapter 8 gives one well-defined version of this formalism, Well-formed Colored PNs, which allows the development of efficient analysis techniques.

Chapters 9, 10 and 11 introduce distributions, which take into account probabilistic properties of systems. They introduce stochastic PNs, or SPNs, and define their semantics in terms of stochastic processes, and, for some classes of models, their relationships with Markov chains. The principal methods of analyzing SPNs are then presented. Chapter 9 introduces stochastic PNs.

Chapter 10 introduces well-formed SPNs by combining the formalisms presented in Chapter 7 (Well-formed Colored PNs) and Chapter 9. Modeling a multiprocessor architecture illustrates the expressivity of this formalism and its interest for performance evaluation. Chapter 11 develops a tensorial composition of classical and well-formed SPNs, showing that such a compositional approach reduces the complexity of the corresponding validation.

Part 2

The second part of this volume presents important advanced analysis techniques and finally gives some significant and illustrative case studies.

Chapters 12 and 13 address checking and verifying non-temporal behaviors. They present the main approaches that are based on building and manipulating the

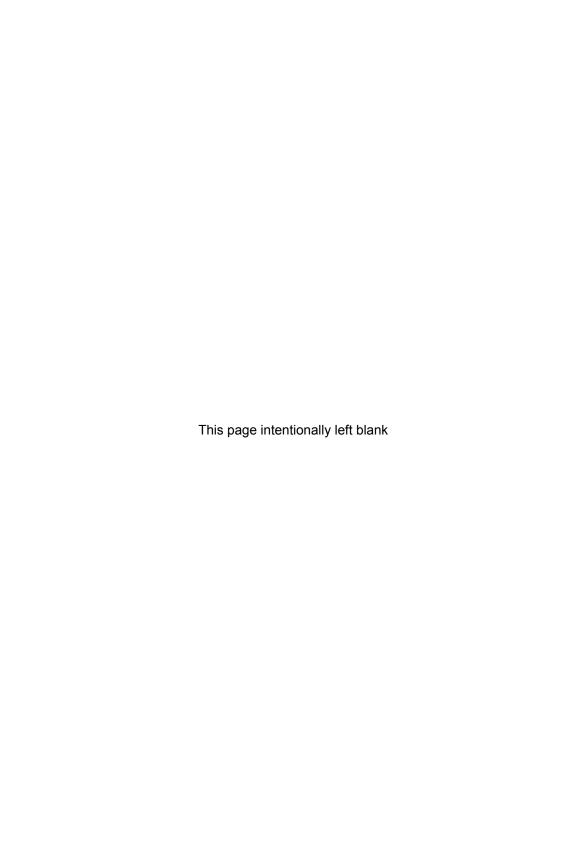
(system) accessibility or reachability graph, i.e. the graph representing all possible behaviors of the model. Checking these properties, by algorithms applied to the accessibility graph, suffers from the problems of combinatorial explosion. The general problems to be solved to control such an increase in the number of states, as well as general solutions, are then given. Three specific techniques, based respectively on the unfolding of the colored PNs, on symmetries, and on partial orders, are then presented.

Chapters 14 and 15 focus on the temporal validation of behaviors. Chapter 14 analyzes the relationships existing between symmetry and temporal logic for the verification of properties that depend on the specificities of the system. Chapter 15 introduces a parallel-serial hierarchy of temporal behaviors. This hierarchy simplifies the description of complex systems and is very well adapted for modeling complex multimedia and hypermedia objects, documents, and systems.

Chapter 16 presents how to use the main relationships that exist between linear logic and Petri nets for specification and validation. Logical reasoning is constructed based on the behavior of PNs that does not need to produce the reachability graph. The interest of linear logic is illustrated by showing in particular how to handle symbolic temporal intervals (minimum, maximum).

Chapters 17 and 18 present two important case studies that are illustrative while being manageable and easily understandable. Chapter 17 is devoted to the modeling and design of a multilayered, multimedia architecture that is able to guarantee temporal properties at the application level. Chapter 18 presents the application of PN-based models to performance evaluation in the field of computer-integrated manufacturing systems.

Finally, a conclusion summarizes the contents of this volume.



Part 1 Fundamental Models

