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Lavinia Corina Ciungu

Non- commutative Multiple-Valued Logic Algebras

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Dedicated to my precious son David Edward

Introduction

In 1920 Łukasiewicz introduced his three valued logic ([223]), the first model of multiple-valued logic. The n -valued propositional logic for $n > 3$ was constructed in 1922 and the \aleph_0 -valued Łukasiewicz-Tarski logic in 1930 ([224]). The first completeness theorem for \aleph_0 -valued Łukasiewicz-Tarski logic was given by Wajsberg in 1935. As a direct generalization of two-valued calculus, Post introduced in 1921 an n -valued propositional calculus distinct from that of Łukasiewicz ([239]).

In the early 1940s Gr.C. Moisil was the first to develop the theory of n -valued Łukasiewicz algebras with the intention of algebraizing Łukasiewicz's logic ([226, 227]), but an example of A. Rose from 1956 established that for $n \geq 5$ the Łukasiewicz implication can no longer be defined on a Łukasiewicz algebra. Consequently, the structures introduced by Moisil are models for Łukasiewicz logic only for $n = 3$ and $n = 4$. These algebras are now called *Łukasiewicz-Moisil algebras* or *LM algebras* for short ([14]).

The loss of implication has led to another type of logic, today called *Moisil logic*, distinct from the Łukasiewicz system. The logic corresponding to n -valued Łukasiewicz-Moisil algebras was created by Moisil in 1964. The fundamental concept of Moisil logic is *nuancing*. During 1954–1973 Moisil introduced the θ -valued LM algebras without negation, applied multiple-valued logics to switching theory and studied algebraic properties of LM algebras (representation, ideals, residuation) ([228]). Moisil's works have been continued by many mathematicians ([149, 151]). A. Iorgulescu introduced and studied θ -valued LM algebras with negation ([170]), while V. Boicescu defined and investigated n -valued LM algebras without negation ([13]).

Today these multiple-valued logics have been developed into fuzzy logics, which connect quantum mechanics, mathematical logic, probability theory, algebra and soft computing.

In 1958 Chang defined *MV-algebras* ([38]) as the algebraic counterpart of \aleph_0 -valued Łukasiewicz logic and he gave another completeness proof of this logic ([39]).

An *MV-algebra* is an algebra $(A, \oplus, \bar{}, 0)$ with a binary operation \oplus , a unary operation $\bar{}$ and a constant 0 satisfying the following equations:

- (MV_1) $(x \oplus y) \oplus z = x \oplus (y \oplus z)$;
 (MV_2) $x \oplus y = y \oplus x$;
 (MV_3) $x \oplus 0 = x$;
 (MV_4) $(x^-)^- = x$;
 (MV_5) $x \oplus 0^- = 0^-$;
 (MV_6) $(x^- \oplus y)^- \oplus y = (y^- \oplus x)^- \oplus x$.

Studies on MV-algebras have been developed in [5–8, 22, 77, 81, 87, 89, 91, 120, 139, 146, 147, 153, 213, 214, 217–219, 247].

Starting from the systems of positive implicational calculus, weak systems of positive implicational calculus and BCI and BCK systems, in 1966 Y. Imai and K. Iséki introduced the BCK-algebras ([168]).

In 1977 R. Grigolia introduced MV_n -algebras to model the n -valued Łukasiewicz logic ([157]) and it was proved that there is a connection between n -valued Łukasiewicz algebras and MV_n -algebras ([171–173, 191, 216]).

One of the most famous results in the theory of MV-algebras was Mundici's theorem from 1986 which states that the category of MV-algebras is equivalent to the category of Abelian ℓ -groups with strong unit ([229]).

The non-commutative generalizations of MV-algebras called *pseudo-MV algebras* were introduced by G. Georgescu and A. Iorgulescu in [135] and [137] and they can be regarded as algebraic semantics for a non-commutative generalization of a multiple-valued reasoning ([215]). The pseudo-MV algebras were introduced independently by J. Rachůnek ([241]) under the name of *generalized MV-algebras*.

A. Dvurečenskij proved in [97] that any pseudo-MV algebra is isomorphic with some interval in an ℓ -group with strong unit, that is, the category of pseudo-MV algebras is equivalent to the category of unital ℓ -groups.

Residuation is a fundamental concept of ordered structures and categories and Ward and Dilworth were the first to introduce the concept of a *residuated lattice* as a generalization of ideal lattices of rings ([262]). The theory of residuated lattices was used to develop algebraic counterparts of fuzzy logics ([256]) and substructural logics ([234]).

A residuated lattice is defined as an algebra $\mathcal{A} = (A, \wedge, \vee, \odot, \rightarrow, \rightsquigarrow, e)$ of type $(2, 2, 2, 2, 2, 0)$ satisfying the following conditions:

- (A_1) (A, \wedge, \vee) is a lattice;
 (A_2) (A, \odot, e) is a monoid;
 (A_3) $x \odot y \leq z$ iff $x \leq y \rightarrow z$ iff $y \leq x \rightsquigarrow z$ for any $x, y, z \in A$ (*pseudo-residuation*).

A residuated lattice with a constant 0 (which can denote any element) is called a *pointed residuated lattice* or *full Lambek algebra* (*FL-algebra*, for short). If $x \leq e$ for all $x \in A$, then \mathcal{A} is called an *integral residuated lattice*. An FL-algebra \mathcal{A} which satisfies the condition $0 \leq x \leq e$ for all $x \in A$ is called *FL_w-algebra* or *bounded integral residuated lattice* ([129]). In this case we put $e = 1$, so that an FL_w-algebra will be denoted $(A, \wedge, \vee, \odot, \rightarrow, \rightsquigarrow, 0, 1)$. Clearly, if \mathcal{A} is an FL_w-algebra, then $(A, \wedge, \vee, 0, 1)$ is a bounded lattice.

In order to formalize the multiple-valued logics induced by continuous t-norms on the real unit interval $[0, 1]$, P. Hájek introduced in 1998 a very general multiple-

valued logic, called *Basic Logic* (or BL) ([158]). Basic Logic turns out to be a common ingredient in three important multiple-valued logics: \aleph_0 -valued Łukasiewicz logic, Gödel logic and Product logic. The Lindenbaum-Tarski algebras for Basic Logic are called *BL-algebras* ([23, 82, 220–222, 255–257]). Apart from their logical interest, BL-algebras have important algebraic properties and they have been intensively studied from an algebraic point of view.

The well-known result that a t-norm on $[0, 1]$ has residuum if and only if the t-norm is left-continuous makes clear that BL is not the most general t-norm based logic. In fact, a weaker logic than BL, called *Monoidal t-norm based logic* (MTL, for short) was defined in [117] and proved in [197] to be the logic of left-continuous t-norms and their residua. The algebraic counterpart of this logic is MTL-algebra, also introduced in [117].

G. Georgescu and A. Iorgulescu introduced in [136] the *pseudo-BL algebras* as a natural generalization of BL-algebras in the non-commutative case. A pseudo-BL algebra is an FL_w -algebra which satisfies the conditions:

$$(A_4) \quad (x \rightarrow y) \odot x = x \odot (x \rightsquigarrow y) = x \wedge y \text{ (pseudo-divisibility);}$$

$$(A_5) \quad (x \rightarrow y) \vee (y \rightarrow x) = (x \rightsquigarrow y) \vee (y \rightsquigarrow x) = 1 \text{ (pseudo-prelinearity).}$$

Properties of pseudo-BL algebras were deeply investigated by A. Di Nola, G. Georgescu and A. Iorgulescu in [85] and [86]. Some classes of pseudo-BL algebras were investigated in [143] and the corresponding propositional logic was established by Hájek in [158] and [159].

A more general structure than the pseudo-BL algebra is the *weak pseudo-BL algebra* or *pseudo-MTL algebra* introduced by P. Flondor, G. Georgescu and A. Iorgulescu in [122]. Pseudo-MTL algebras are FL_w -algebras satisfying condition (A_5) and they include as a particular case the *weak BL-algebras* which is an alternative name for MTL-algebras.

Properties of pseudo-MTL algebras are also studied in [46, 144, 181].

An FL_w -algebra which satisfies condition (A_4) is called a *divisible residuated lattice* or *bounded $R\ell$ -monoid*. Properties of divisible residuated lattices were studied by A. Dvurečenskij, J. Rachůnek and J. Kühr ([105, 111, 205, 240]).

Pseudo-BCK algebras were introduced in 2001 by G. Georgescu and A. Iorgulescu ([138]) as non-commutative generalizations of BCK-algebras. Properties of pseudo-BCK algebras and their connection with other fuzzy structures were established by A. Iorgulescu in [179–182].

For a guide through the pseudo-BCK algebras realm we refer the reader to the monograph [186].

Another generalization of pseudo-BL algebras was given in [148], where *pseudo-hoops* were defined and studied. Pseudo-hoops were originally introduced by Bosbach in [15] and [16] under the name of *complementary semigroups*. It was proved that a pseudo-hoop has the pseudo-divisibility condition and it is a meet-semilattice, so a bounded $R\ell$ -monoid can be viewed as a bounded pseudo-hoop together with the join-semilattice property. In other words, a bounded pseudo-hoop is a meet-semilattice ordered residuated, integral and divisible monoid.

Other topics in multiple-valued logic algebras have been studied in [34, 36, 92, 132, 141, 150, 248].

The notion of a *state* is an analogue of a probability measure and it has a very important role in the theory of quantum structures ([108]). The basic idea of states is an averaging of events (elements) of a given algebraic structure. Since in the case of Łukasiewicz ∞ -valued logic the set of events has the structure of an MV-algebra, the theory of probability on this logic is based on the notion of a state defined on an MV-algebra. Besides mathematical logic, Riečan and Neubrunn studied MV-algebras as fields of events in generalized probability theory ([250]). Therefore, the study of states on MV-algebras is a very active field of research ([40, 83, 84, 119, 133, 246]) which arises from the general problem of investigating probabilities defined for logical systems.

States on an MV-algebra $(A, \oplus, \bar{\cdot}, 0)$ were first introduced by D. Mundici in [230] as functions $s : A \rightarrow [0, 1]$ satisfying the conditions:

$$\begin{aligned} s(1) &= 1 \text{ (normality);} \\ s(x \oplus y) &= s(x) + s(y) \text{ if } x \odot y = 0 \text{ (additivity),} \end{aligned}$$

where $x \odot y = (x^- \oplus y^-)^-$.

They are analogous to finitely additive probability measures on Boolean algebras and play a crucial role in MV-algebraic probability theory ([249]).

States on other commutative and non-commutative algebraic structures have been defined and investigated by many authors ([20, 21, 102, 133, 134, 140, 142, 258, 259]).

The aim of this book is to present new results regarding non-commutative multiple-valued logic algebras and some of their applications. Almost all the results are based on the author's recent papers ([42–75]).

The book consists of nine chapters.

The Chap. 1 is devoted to pseudo-BCK algebras. After presenting the basic definitions and properties, we prove new properties of pseudo-BCK algebras with pseudo-product and pseudo-BCK algebras with pseudo-double negation. Examples of proper pseudo-BCK algebras, good pseudo-BCK algebras and pseudo-BCK lattices are given, and the orthogonal elements in a pseudo-BCK algebra are characterized. Finally, we define the maximal and normal deductive systems of a pseudo-BCK algebra with pseudo-product and we study their properties.

In Chap. 2 we recall the basic properties of pseudo-hoops, we introduce the notions of join-center and cancellative-center of pseudo-hoops and we define and study algebras on subintervals of pseudo-hoops. Additionally, new properties of a pseudo-hoop are proved.

Chapter 3 is devoted to residuated lattices. We investigate the properties of the Boolean center of an FL_w -algebra and we define and study the directly indecomposable FL_w -algebras. One of the main results consists of proving that any linearly ordered FL_w -algebra is directly indecomposable. Finally, we define and study FL_w -algebras of fractions relative to a meet-closed system.

In Chap. 4 we present some specific properties of other non-commutative multiple-valued logic algebras: pseudo-MTL algebras, bounded $R\ell$ -monoids, pseudo-BL algebras and pseudo-MV algebras. As main results, we extend to the case of pseudo-MTL algebras some results regarding prime filters proved for

pseudo-BL algebras. The Glivenko property for a good pseudo-BCK algebra is defined and it is shown that a good pseudo-hoop has the Glivenko property.

Chapter 5 deals with special classes of non-commutative residuated structures: local, perfect and Archimedean structures. The local bounded pseudo-BCK(pP) algebras are characterized in terms of primary deductive systems, while the perfect pseudo-BCK(pP) algebras are characterized in terms of perfect deductive systems. One of the main results consists of proving that the radical of a bounded pseudo-BCK(pP) algebra is a normal deductive system. We also prove that any linearly ordered pseudo-BCK(pP) algebra and any locally finite pseudo-BCK(pP) algebra are local. Other results state that any local FL_w -algebra and any locally finite FL_w -algebra are directly indecomposable. The classes of Archimedean and hyperarchimedean FL_w -algebras are introduced and it is proved that any locally finite FL_w -algebra is hyperarchimedean and any hyperarchimedean FL_w -algebra is Archimedean.

Chapter 6 is devoted to the presentation of states on multiple-valued logic algebras. We introduce the notion of states on pseudo-BCK algebras and we study their properties. One of the main results consists of proving that any Bosbach state on a good pseudo-BCK algebra is a Riečan state, however the converse turns out not to be true. We also prove that every Riečan state on a good pseudo-BCK algebra with pseudo-double negation is a Bosbach state. In contrast to the case of pseudo-BL algebras, we show that there exist linearly ordered pseudo-BCK algebras having no Bosbach states and that there exist pseudo-BCK algebras having normal filters which are maximal, but having no Bosbach states.

Some specific properties of states on FL_w -algebras, pseudo-MTL algebras, bounded $R\ell$ -monoids and subinterval algebras of pseudo-hoops are proved.

A special section is dedicated to the existence of states on the residuated structures, showing that every perfect FL_w -algebra admits at least a Bosbach state and every perfect pseudo-BL algebra has a unique state-morphism.

Finally, we introduce the notion of a local state on a perfect pseudo-MTL algebra and we prove that every local state can be extended to a Riečan state.

In Chap. 7 we generalize measures on BCK algebras introduced by A. Dvurečenskij in [94] and [108] to pseudo-BCK algebras that are not necessarily bounded. In particular, we show that if A is a downwards-directed pseudo-BCK algebra and m a measure on it, then the quotient over the kernel of m can be embedded into the negative cone of an Abelian, Archimedean ℓ -group as its subalgebra. This result will enable us to characterize nonzero measure-morphisms on downwards-directed pseudo-BCK algebras as measures whose kernel is a maximal filter. We study state-measures on pseudo-BCK algebras with strong unit and we show how to characterize state-measure-morphisms as extremal state-measures or as state-measures whose kernel is a maximal filter. In particular, we show that for unital pseudo-BCK algebras that are downwards-directed, the quotient over the kernel can be embedded into the negative cone of an Abelian, Archimedean ℓ -group with strong unit. We generalize to pseudo-BCK algebras the identity between de Finetti maps and Bosbach states, following the results proved by Kühr and Mundici in [211] who showed that de Finetti's coherence principle, which has its origins in Dutch bookmaking, has

a strong relationship with MV-states on MV-algebras. We also generalize this for state-measures on unital pseudo-BCK algebras that are downwards-directed.

Chapter 8 is devoted to generalized states on residuated structures. The study of these generalized states is motivated by their interpretation as a new type of semantics for non-commutative fuzzy logics. Usually, the truth degree of sentences in a fuzzy logic is a number in the interval $[0, 1]$ or, more generally, an element of an FL_w -algebra. Similarly, for generalized states, the probability of sentences is evaluated in an arbitrary FL_w -algebra.

We define the generalized states of type I and type II and generalized state-morphisms and we study the relationship between them. We prove that any perfect FL_w -algebra admits strong type I and type II states. Some conditions are given for a generalized state of type I on a linearly ordered bounded $R\ell$ -monoid to be a state operator. The notion of a strong perfect FL_w -algebra is introduced and it is proved that any strong perfect FL_w -algebra admits a generalized state-morphism. The notion of a generalized Riečan state is also introduced and the main results are proved based on the Glivenko property defined for the non-commutative case. The main results consist of proving that any order-preserving type I state is a generalized Riečan state and in some particular conditions the two states coincide. We introduce the notion of a generalized local state on a perfect pseudo-MTL algebra A and we prove that, if A is relatively free of zero divisors, then every generalized local state can be extended to a generalized Riečan state.

Chapter 9 deals with residuated structures with internal states. We define the notions of state operator, strong state operator, state-morphism operator, weak state-morphism operator and we study their properties. We prove that every strong state pseudo-hoop is a state pseudo-hoop and any state operator on an idempotent pseudo-hoop is a weak state-morphism operator. It is proved that for an idempotent pseudo-hoop A a state operator on $\text{Reg}(A)$ can be extended to a state operator on A . One of the main results of this chapter consists of proving that every perfect pseudo-hoop admits a nontrivial state operator. Other results compare the state operators with states and generalized states on a pseudo-hoop. Some conditions are given for a state operator to be a generalized state and for a generalized state to be a state operator.

We hope that this book will be useful to graduate students and researchers in the area of algebras of multiple-valued logics.

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Chapter 1

Pseudo-BCK Algebras

BCK algebras were originally introduced by K. Isèki in [194] with a binary operation $*$ modeling the set-theoretical difference and with a constant element 0 , that is, a least element. Another motivation is from classical and non-classical propositional calculi modeling logical implications. Such algebras contain as a special subfamily the family of MV-algebras where some important fuzzy structures can be studied. For more about BCK algebras, see [167, 174–179, 182–187, 189, 192, 193, 225].

Pseudo-BCK algebras were introduced by G. Georgescu and A. Iorgulescu in [138] as algebras with “two differences”, a left- and right-difference, instead of one $*$ and with a constant element 0 as the least element. In [112], a special subclass of pseudo-BCK algebras, called Łukasiewicz pseudo-BCK algebras, was introduced and it was shown that each such algebra is always a subalgebra of the positive cone of some ℓ -group (not necessarily Abelian). The class of Łukasiewicz pseudo-BCK algebras is a variety whereas the class of pseudo-BCK algebras is not; it is only a quasivariety because it is not closed under homomorphic images. Nowadays pseudo-BCK algebras are used in a dual form, with two implications, \rightarrow and \rightsquigarrow and with one constant element 1 , that is the greatest element. Thus such pseudo-BCK algebras are in the “negative cone” and are also called “left-ones”. Further properties of pseudo-BCK algebras and their connection with other fuzzy structures were established by A. Iorgulescu in [179–182]. For a guide through the pseudo-BCK algebras realm, see the monograph [186]. Studies on pseudo-BCK algebras were also developed in [107, 163, 190, 206, 208–210].

In this chapter we prove new properties of pseudo-BCK algebras with pseudo-product and pseudo-BCK algebras with pseudo-double negation and we show that every pseudo-BCK algebra can be extended to a good one. Examples of proper pseudo-BCK algebras, good pseudo-BCK algebras and pseudo-BCK lattices are given and the orthogonal elements in a pseudo-BCK algebra are characterized. Finally, we define the maximal and normal deductive systems of a pseudo-BCK algebra with pseudo-product and we study their properties.

1.1 Definitions and Properties

Definition 1.1 A *pseudo-BCK algebra* (more precisely, *reversed left-pseudo-BCK algebra*) is a structure $\mathcal{A} = (A, \leq, \rightarrow, \rightsquigarrow, 1)$ where \leq is a binary relation on A , \rightarrow and \rightsquigarrow are binary operations on A and 1 is an element of A satisfying, for all $x, y, z \in A$, the axioms:

$$(psBCK_1) \quad x \rightarrow y \leq (y \rightarrow z) \rightsquigarrow (x \rightarrow z), x \rightsquigarrow y \leq (y \rightsquigarrow z) \rightarrow (x \rightsquigarrow z);$$

$$(psBCK_2) \quad x \leq (x \rightarrow y) \rightsquigarrow y, x \leq (x \rightsquigarrow y) \rightarrow y;$$

$$(psBCK_3) \quad x \leq x;$$

$$(psBCK_4) \quad x \leq 1;$$

$$(psBCK_5) \quad \text{if } x \leq y \text{ and } y \leq x, \text{ then } x = y;$$

$$(psBCK_6) \quad x \leq y \text{ iff } x \rightarrow y = 1 \text{ iff } x \rightsquigarrow y = 1.$$

A pseudo-BCK algebra $\mathcal{A} = (A, \leq, \rightarrow, \rightsquigarrow, 1)$ is *commutative* if $\rightarrow = \rightsquigarrow$. Any commutative pseudo-BCK algebra is a BCK-algebra.

In the sequel we will refer to the pseudo-BCK algebra $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ by its universe A .

Proposition 1.1 *The structure $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK algebra iff the algebra $(A, \rightarrow, \rightsquigarrow, 1)$ of type $(2, 2, 0)$ satisfies the following identities and quasi-identity:*

$$(psBCK'_1) \quad (x \rightarrow y) \rightsquigarrow [(y \rightarrow z) \rightsquigarrow (x \rightarrow z)] = 1;$$

$$(psBCK'_2) \quad (x \rightsquigarrow y) \rightarrow [(y \rightsquigarrow z) \rightarrow (x \rightsquigarrow z)] = 1;$$

$$(psBCK'_3) \quad 1 \rightarrow x = x;$$

$$(psBCK'_4) \quad 1 \rightsquigarrow x = x;$$

$$(psBCK'_5) \quad x \rightarrow 1 = 1;$$

$$(psBCK'_6) \quad (x \rightarrow y = 1 \text{ and } y \rightarrow x = 1) \text{ implies } x = y.$$

Proof Obviously, any pseudo-BCK algebra satisfies $(psBCK'_1)$ – $(psBCK'_6)$.

Conversely, assume that an algebra $(A, \rightarrow, \rightsquigarrow, 1)$ satisfies $(psBCK'_1)$ – $(psBCK'_6)$.

Applying $(psBCK'_3)$ and $(psBCK'_1)$ we get:

$$x \rightsquigarrow [(x \rightarrow y) \rightsquigarrow y] = (1 \rightarrow x) \rightsquigarrow [(x \rightarrow y) \rightsquigarrow (1 \rightarrow y)] = 1.$$

Similarly, by $(psBCK'_4)$ and $(psBCK'_2)$ we have:

$$x \rightarrow [(x \rightsquigarrow y) \rightarrow y] = (1 \rightsquigarrow x) \rightarrow [(x \rightsquigarrow y) \rightarrow (1 \rightsquigarrow y)] = 1.$$

Applying $(psBCK'_3)$ and $(psBCK'_2)$ we have:

$$x \rightarrow x = 1 \rightarrow (x \rightarrow x) = (1 \rightsquigarrow 1) \rightarrow [(1 \rightsquigarrow x) \rightarrow (1 \rightsquigarrow x)] = 1.$$

Similarly, by $(psBCK'_4)$ and $(psBCK'_1)$ we get:

$$x \rightsquigarrow x = 1 \rightsquigarrow (x \rightsquigarrow x) = (1 \rightarrow 1) \rightsquigarrow [(1 \rightarrow x) \rightsquigarrow (1 \rightarrow x)] = 1.$$

Moreover, if $x \rightarrow y = 1$ then $x \rightsquigarrow y = x \rightsquigarrow [(x \rightarrow y) \rightsquigarrow y] = 1$ and similarly, if $x \rightsquigarrow y = 1$ then $x \rightarrow y = x \rightarrow [(x \rightsquigarrow y) \rightarrow y] = 1$.

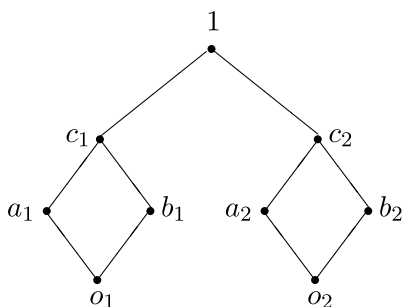


Fig. 1.1 Example of proper pseudo-BCK algebra

It follows that $x \rightarrow y = 1$ iff $x \rightsquigarrow y = 1$.

We deduce that the relation \leq defined by $x \leq y$ iff $x \rightarrow y = 1$ is a partial order on A which makes $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ a pseudo-BCK algebra. \square

In the sequel, we shall use either $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ or $(A, \rightarrow, \rightsquigarrow, 1)$ for a pseudo-BCK algebra.

Example 1.1 Consider $A = \{o_1, a_1, b_1, c_1, o_2, a_2, b_2, c_2, 1\}$ with $o_1 < a_1, b_1 < c_1 < 1$ and a_1, b_1 incomparable, $o_2 < a_2, b_2 < c_2 < 1$ and a_2, b_2 incomparable. Assume that any element of the set $\{o_1, a_1, b_1, c_1\}$ is incomparable with any element of the set $\{o_2, a_2, b_2, c_2\}$ (see Fig. 1.1).

Consider the operations $\rightarrow, \rightsquigarrow$ given by the following tables:

\rightarrow	o_1	a_1	b_1	c_1	o_2	a_2	b_2	c_2	1
o_1	1	1	1	1	o_2	a_2	b_2	c_2	1
a_1	o_1	1	b_1	1	o_2	a_2	b_2	c_2	1
b_1	a_1	a_1	1	1	o_2	a_2	b_2	c_2	1
c_1	o_1	a_1	b_1	1	o_2	a_2	b_2	c_2	1
o_2	o_1	a_1	b_1	c_1	1	1	1	1	1
a_2	o_1	a_1	b_1	c_1	o_2	1	b_2	1	1
b_2	o_1	a_1	b_1	c_1	c_2	c_2	1	1	1
c_2	o_1	a_1	b_1	c_1	o_2	c_2	b_2	1	1
1	o_1	a_1	b_1	c_1	o_2	a_2	b_2	c_2	1
\rightsquigarrow	o_1	a_1	b_1	c_1	o_2	a_2	b_2	c_2	1
o_1	1	1	1	1	o_2	a_2	b_2	c_2	1
a_1	b_1	1	b_1	1	o_2	a_2	b_2	c_2	1
b_1	o_1	a_1	1	1	o_2	a_2	b_2	c_2	1
c_1	o_1	a_1	b_1	1	o_2	a_2	b_2	c_2	1
o_2	o_1	a_1	b_1	c_1	1	1	1	1	1
a_2	o_1	a_1	b_1	c_1	b_2	1	b_2	1	1
b_2	o_1	a_1	b_1	c_1	b_2	c_2	1	1	1
c_2	o_1	a_1	b_1	c_1	b_2	c_2	b_2	1	1
1	o_1	a_1	b_1	c_1	o_2	a_2	b_2	c_2	1

Then $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ is a proper pseudo-BCK algebra.

We recall the definition of an ℓ -group. The language of lattice-ordered groups (ℓ -groups) involves both the group operations and the binary lattice operations.

By a *lattice-ordered group* (ℓ -group) we will mean an ordered group (G, \leq) such that (G, \leq) is a lattice. The ℓ -group G is called an ℓu -group if there exists an element $u > 0$ such that for any $x \in G$ there is an $n \in \mathbb{N}$ such that $x \leq nu$. The element u is called a *strong unit*.

For details regarding ℓ -groups we refer the reader to [2, 12, 76].

Example 1.2 Let $(G, \vee, \wedge, +, -, 0)$ be an ℓ -group.

On the negative cone $G^- = \{g \in G \mid g \leq 0\}$ we define:

$$\begin{aligned} g \rightarrow h &:= h - (g \vee h) = (h - g) \wedge 0, \\ g \rightsquigarrow h &:= -(g \vee h) + h = (-g + h) \wedge 0. \end{aligned}$$

Then $(G^-, \leq, \rightarrow, \rightsquigarrow, 0)$ is a pseudo-BCK algebra.

Remark 1.1 (Definition of union) Let $(A_i, \leq, \rightarrow_i, \rightsquigarrow_i, 1_i)_{i \in I}$ be a collection of pseudo-BCK algebras such that:

- (i) $1_i = 1$ for all $i \in I$,
- (ii) $A_i \cap A_j = \{1\}$ for all $i, j \in I, i \neq j$.

Let $A = \bigcup_{i \in I} A_i$ and define:

$$\begin{aligned} x \rightarrow y &:= \begin{cases} x \rightarrow_i y & \text{if } x, y \in A_i, i \in I \\ y & \text{otherwise,} \end{cases} \\ x \rightsquigarrow y &:= \begin{cases} x \rightsquigarrow_i y & \text{if } x, y \in A_i, i \in I \\ y & \text{otherwise.} \end{cases} \end{aligned}$$

Then $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK algebra called the *union* of the pseudo-BCK algebras $(A_i, \leq, \rightarrow_i, \rightsquigarrow_i, 1_i)_{i \in I}$.

Note that the notion of union defined above is not related to the notion of ordinal sum defined in Chap. 2.

Proposition 1.2 *In any pseudo-BCK algebra A the following properties hold:*

- (psbck-c₁) $x \leq y$ implies $y \rightarrow z \leq x \rightarrow z$ and $y \rightsquigarrow z \leq x \rightsquigarrow z$;
- (psbck-c₂) $x \leq y, y \leq z$ implies $x \leq z$;
- (psbck-c₃) $x \rightarrow (y \rightsquigarrow z) = y \rightsquigarrow (x \rightarrow z), x \rightsquigarrow (y \rightarrow z) = y \rightarrow (x \rightsquigarrow z)$;
- (psbck-c₄) $z \leq y \rightarrow x$ iff $y \leq z \rightsquigarrow x$;
- (psbck-c₅) $z \rightarrow x \leq (y \rightarrow z) \rightarrow (y \rightarrow x), z \rightsquigarrow x \leq (y \rightsquigarrow z) \rightsquigarrow (y \rightsquigarrow x)$;
- (psbck-c₆) $x \leq y \rightarrow x, x \leq y \rightsquigarrow x$;

(*psbck-c7*) $1 \rightarrow x = x = 1 \rightsquigarrow x$;

(*psbck-c8*) $x \rightarrow x = x \rightsquigarrow x = 1$;

(*psbck-c9*) $x \rightarrow 1 = x \rightsquigarrow 1 = 1$;

(*psbck-c10*) $x \leq y$ implies $z \rightarrow x \leq z \rightarrow y$ and $z \rightsquigarrow x \leq z \rightsquigarrow y$;

(*psbck-c11*) $[(y \rightarrow x) \rightsquigarrow x] \rightarrow x = y \rightarrow x$, $[(y \rightsquigarrow x) \rightarrow x] \rightsquigarrow x = y \rightsquigarrow x$.

Proof

(*psbck-c1*) Since $x \leq y$, applying (*psBCK₆*), (*psBCK₁*) and (*psBCK₄*) we get $1 = x \rightarrow y \leq (y \rightarrow z) \rightsquigarrow (x \rightarrow z)$, so $(y \rightarrow z) \rightsquigarrow (x \rightarrow z) = 1$ for all $z \in A$.

Applying (*psBCK₆*) again we get $y \rightarrow z \leq x \rightarrow z$.

Similarly, $y \rightsquigarrow z \leq x \rightsquigarrow z$.

(*psbck-c2*) By (*psbck-c1*), $x \leq y$ implies $y \rightarrow z \leq x \rightarrow z$. Since $y \leq z$ we have $y \rightarrow z = 1$, so $x \rightarrow z = 1$. Applying (*psBCK₆*) we get $x \leq z$.

(*psbck-c3*) Applying (*psBCK₁*) we have $y \rightarrow x \leq (x \rightarrow z) \rightsquigarrow (y \rightarrow z)$ and by (*psbck-c1*) we get $[(x \rightarrow z) \rightsquigarrow (y \rightarrow z)] \rightsquigarrow u \leq (y \rightarrow x) \rightsquigarrow u$ for any $u \in A$.

From this inequality, replacing z with $u \rightsquigarrow z$, x with $x \rightsquigarrow z$ and u with $(u \rightsquigarrow x) \rightsquigarrow [y \rightarrow (u \rightsquigarrow z)]$ we get

$$\begin{aligned} & [[(x \rightsquigarrow z) \rightarrow (u \rightsquigarrow z)] \rightsquigarrow [y \rightarrow (u \rightsquigarrow z)]] \rightsquigarrow [(u \rightsquigarrow x) \rightsquigarrow [y \rightarrow (u \rightsquigarrow z)]] \\ & \leq [y \rightarrow (x \rightsquigarrow z)] \rightsquigarrow [(u \rightsquigarrow x) \rightsquigarrow [y \rightarrow (u \rightsquigarrow z)]]. \end{aligned}$$

By (*psBCK₁*) we have $u \rightsquigarrow x \leq (x \rightsquigarrow z) \rightarrow (u \rightsquigarrow z)$ and applying (*psbck-c1*) it follows that the left-hand side of the above inequality is equal to 1.

Thus the right-hand side is also equal to 1, so $y \rightarrow (x \rightsquigarrow z) \leq (u \rightsquigarrow x) \rightsquigarrow [y \rightarrow (u \rightsquigarrow z)]$.

Replacing x with $y \rightarrow z$ and u with x we get

$$y \rightarrow [(y \rightarrow z) \rightsquigarrow z] \leq [x \rightsquigarrow (y \rightarrow z)] \rightsquigarrow [y \rightarrow (x \rightsquigarrow z)].$$

But, by (*psBCK₂*) we have $y \leq (y \rightarrow z) \rightsquigarrow z$, so $y \rightarrow [(y \rightarrow z) \rightsquigarrow z] = 1$.

It follows that $[x \rightsquigarrow (y \rightarrow z)] \rightsquigarrow [y \rightarrow (x \rightsquigarrow z)] = 1$.

Therefore $x \rightsquigarrow (y \rightarrow z) \leq y \rightarrow (x \rightsquigarrow z)$.

On the other hand, by (*psBCK₂*) we have $x \leq (x \rightsquigarrow z) \rightarrow z$ and applying (*psbck-c1*) we get $[(x \rightsquigarrow z) \rightarrow z] \rightsquigarrow (y \rightarrow z) \leq x \rightsquigarrow (y \rightarrow z)$.

By (*psBCK₁*) we have $y \rightarrow x \leq (x \rightarrow z) \rightsquigarrow (y \rightarrow z)$ and replacing x with $x \rightsquigarrow z$ we get $y \rightarrow (x \rightsquigarrow z) \leq [(x \rightsquigarrow z) \rightarrow z] \rightsquigarrow (y \rightarrow z) \leq x \rightsquigarrow (y \rightarrow z)$.

We conclude that $x \rightarrow (y \rightsquigarrow z) = y \rightsquigarrow (x \rightarrow z)$.

Similarly, $x \rightsquigarrow (y \rightarrow z) = y \rightarrow (x \rightsquigarrow z)$.

(*psbck-c4*) From $z \leq y \rightarrow x$, by (*psBCK₂*) and (*psbck-c1*) we have

$$y \leq (y \rightarrow x) \rightsquigarrow x \leq z \rightsquigarrow x.$$

Similarly, from $y \leq z \rightsquigarrow x$ we get $z \leq (z \rightsquigarrow x) \rightarrow x \leq y \rightarrow x$.

(*psbck-c5*) Applying (*psBCK₁*) we have $y \rightarrow z \leq (z \rightarrow x) \rightsquigarrow (y \rightarrow x)$ and according to (*psbck-c1*) we get

$$[(z \rightarrow x) \rightsquigarrow (y \rightarrow x)] \rightarrow (y \rightarrow x) \leq (y \rightarrow z) \rightarrow (y \rightarrow x).$$

By (*psBCK₂*) it follows that $z \rightarrow x \leq [(z \rightarrow x) \rightsquigarrow (y \rightarrow x)] \rightarrow (y \rightarrow x)$, and applying (*psbck-c2*) we conclude that $z \rightarrow x \leq (y \rightarrow z) \rightarrow (y \rightarrow x)$.

Similarly, from $y \rightsquigarrow z \leq (z \rightsquigarrow x) \rightarrow (y \rightsquigarrow x)$ we get $z \rightsquigarrow x \leq (y \rightsquigarrow z) \rightsquigarrow (y \rightsquigarrow x)$.

(*psbck-c6*) Since $y \leq 1 = x \rightarrow x$, it follows by (*psbck-c4*) that $x \leq y \rightsquigarrow x$.

Similarly, from $y \leq 1 = x \rightsquigarrow x$ we get $x \leq y \rightarrow x$.

(*psbck-c7*) By (*psbck-c6*) we have $x \leq 1 \rightarrow x$ and $x \leq 1 \rightsquigarrow x$.

By (*psBCK₂*) we get $1 \leq (1 \rightarrow x) \rightsquigarrow x$ and $1 \leq (1 \rightsquigarrow x) \rightarrow x$.

It follows that $(1 \rightarrow x) \rightsquigarrow x = 1$ and $(1 \rightsquigarrow x) \rightarrow x = 1$, so $1 \rightarrow x \leq x$ and $1 \rightsquigarrow x \leq x$. Thus $1 \rightarrow x = x = 1 \rightsquigarrow x$.

(*psbck-c8*) and (*psbck-c9*) are consequences of the axiom (*psBCK₆*).

(*psbck-c10*) Applying (*psbck-c7*), (*psBCK₆*) and (*psBCK₁*) we have:

$$\begin{aligned} z \rightarrow y = 1 \rightsquigarrow (z \rightarrow y) &= (x \rightarrow y) \rightsquigarrow (z \rightarrow y) \geq z \rightarrow x \quad \text{and} \\ z \rightsquigarrow y = 1 \rightarrow (z \rightsquigarrow y) &= (x \rightsquigarrow y) \rightarrow (z \rightsquigarrow y) \geq z \rightsquigarrow x. \end{aligned}$$

(*psbck-c11*) By (*psBCK₂*) we have $y \leq (y \rightarrow x) \rightsquigarrow x$ and $y \leq (y \rightsquigarrow x) \rightarrow x$.

Applying (*psbck-c1*) we get

$$[(y \rightarrow x) \rightsquigarrow x] \rightarrow x \leq y \rightarrow x \quad \text{and} \quad [(y \rightsquigarrow x) \rightarrow x] \rightsquigarrow x \leq y \rightsquigarrow x.$$

On the other hand, by (*psBCK₂*) we have:

$$y \rightarrow x \leq [(y \rightarrow x) \rightsquigarrow x] \rightarrow x \quad \text{and} \quad y \rightsquigarrow x \leq [(y \rightsquigarrow x) \rightarrow x] \rightsquigarrow x.$$

We conclude that

$$[(y \rightarrow x) \rightsquigarrow x] \rightarrow x = y \rightarrow x \quad \text{and} \quad [(y \rightsquigarrow x) \rightarrow x] \rightsquigarrow x = y \rightsquigarrow x. \quad \square$$

Proposition 1.3 *Let $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ be a pseudo-BCK algebra.*

If $\bigvee_{i \in I} x_i$ exists, then so does $\bigwedge_{i \in I} (x_i \rightarrow y)$ and $\bigwedge_{i \in I} (x_i \rightsquigarrow y)$ and we have:

(*psbck-c12*) $(\bigvee_{i \in I} x_i) \rightarrow y = \bigwedge_{i \in I} (x_i \rightarrow y)$, $(\bigvee_{i \in I} x_i) \rightsquigarrow y = \bigwedge_{i \in I} (x_i \rightsquigarrow y)$.

Proof If we let $x = \bigvee_{i \in I} x_i$, it follows that $x_i \leq x$ and applying (*psbck-c1*) we have $x \rightarrow y \leq x_i \rightarrow y$ for all $i \in I$. Let z be a lower bound of $\{x_i \rightarrow y \mid i \in I\}$. Then, by (*psbck-c4*), $z \leq x_i \rightarrow y$ implies $x_i \leq z \rightsquigarrow y$ for all $i \in I$, so $x \leq z \rightsquigarrow y$. Applying (*psbck-c4*) again, we get $z \leq x \rightarrow y$.

Thus $x \rightarrow y$ is the g.l.b. of $\{x_i \rightarrow y \mid i \in I\}$.

We conclude that $\bigwedge_{i \in I} (x_i \rightarrow y)$ exists and $(\bigvee_{i \in I} x_i) \rightarrow y = \bigwedge_{i \in I} (x_i \rightarrow y)$.

Similarly, $\bigwedge_{i \in I} (x_i \rightsquigarrow y)$ exists and $(\bigvee_{i \in I} x_i) \rightsquigarrow y = \bigwedge_{i \in I} (x_i \rightsquigarrow y)$. \square

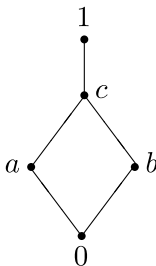


Fig. 1.2 Example of bounded pseudo-BCK algebra

Definition 1.2 If there is an element 0 of a pseudo-BCK algebra $(A, \leq, \rightarrow, \rightsquigarrow, 1)$, such that $0 \leq x$ (i.e. $0 \rightarrow x = 0 \rightsquigarrow x = 1$), for all $x \in A$, then 0 is called the *zero* of A . A pseudo-BCK algebra with zero is called a *bounded pseudo-BCK algebra* and it is denoted by $(A, \leq, \rightarrow, \rightsquigarrow, 0, 1)$.

Example 1.3 Consider $A = \{0, a, b, c, 1\}$ with $0 < a, b < c < 1$ and a, b incomparable (see Fig. 1.2).

Consider the operations $\rightarrow, \rightsquigarrow$ given by the following tables:

\rightarrow	0	a	b	c	1
0	1	1	1	1	1
a	0	1	b	1	1
b	a	a	1	1	1
c	0	a	b	1	1
1	0	a	b	c	1

\rightsquigarrow	0	a	b	c	1
0	1	1	1	1	1
a	b	1	b	1	1
b	0	a	1	1	1
c	0	a	b	1	1
1	0	a	b	c	1

Then $(A, \leq, \rightarrow, \rightsquigarrow, 0, 1)$ is a bounded pseudo-BCK algebra. (As we will see later, A is even a pseudo-BCK lattice.)

Let $(A, \leq, \rightarrow, \rightsquigarrow, 0, 1)$ be a bounded pseudo-BCK algebra. We define two negations $\bar{}$ and $\tilde{}$: for all $x \in A$,

$$x^- := x \rightarrow 0, \quad x^\sim := x \rightsquigarrow 0.$$

In the sequel we will use the following notation:

$$x^{--} = (x^-)^-; \quad x^{\sim\sim} = (x^\sim)^\sim; \quad x^{-\sim} = (x^-)^\sim; \quad x^{\sim-} = (x^\sim)^-.$$

Example 1.4 Let $(G, \vee, \wedge, +, -, 0)$ be an ℓ -group with a strong unit $u \geq 0$. On the interval $[-u, 0]$ we define:

$$x \rightarrow y := (y - x) \wedge 0, \quad x \rightsquigarrow y := (-x + y) \wedge 0.$$

Then $([-u, 0], \leq, \rightarrow, \rightsquigarrow, -u, 0)$ is a bounded pseudo-BCK algebra with $x^- = -u - x$ and $x^\sim = -x - u$. In a similar way, $([-u, 0], \leq, \rightarrow, \rightsquigarrow, 0)$ is a pseudo-BCK algebra that is not bounded.

Example 1.5 Let $(G, \vee, \wedge, +, -, 0)$ be an ℓ -group with a strong unit $u \geq 0$. On the interval $[0, u]$ we define:

$$x \rightarrow y := (u - x + y) \wedge u, \quad x \rightsquigarrow y := (y - x + u) \wedge u.$$

Then $([0, u], \leq, \rightarrow, \rightsquigarrow, 0, u)$ is a bounded pseudo-BCK algebra with $x^- = u - x$ and $x^\sim = -x + u$. If on $[0, u]$ we set $\rightarrow_1 = \rightsquigarrow$ and $\rightsquigarrow_1 = \rightarrow$, then $([0, u], \leq, \rightarrow_1, \rightsquigarrow_1, 0, u)$ is isomorphic with $([-u, 0], \leq, \rightarrow, \rightsquigarrow, -u, 0)$ under the isomorphism $x \mapsto x - u, x \in [0, u]$.

Proposition 1.4 *In a bounded pseudo-BCK algebra the following hold:*

- (psbck-c13) $1^- = 0 = 1^\sim, 0^- = 1 = 0^\sim$;
- (psbck-c14) $x \leq x^{--}, x \leq x^{\sim\sim}$;
- (psbck-c15) $x \rightarrow y \leq y^- \rightsquigarrow x^-, x \rightsquigarrow y \leq y^\sim \rightarrow x^\sim$;
- (psbck-c16) $x \leq y$ implies $y^- \leq x^-$ and $y^\sim \leq x^\sim$;
- (psbck-c17) $x \rightarrow y^\sim = y \rightsquigarrow x^-$ and $x \rightsquigarrow y^- = y \rightarrow x^\sim$;
- (psbck-c18) $x^{--} = x^-, x^{\sim\sim} = x^\sim$;
- (psbck-c19) $x \rightarrow y^{\sim\sim} = y^- \rightsquigarrow x^- = x^{\sim\sim} \rightarrow y^{\sim\sim}$ and $x \rightsquigarrow y^{\sim\sim} = y^\sim \rightarrow x^\sim = x^{\sim\sim} \rightsquigarrow y^{\sim\sim}$;
- (psbck-c20) $x \rightarrow y^\sim = y^{\sim\sim} \rightsquigarrow x^- = x^{\sim\sim} \rightarrow y^\sim$ and $x \rightsquigarrow y^- = y^{\sim\sim} \rightarrow x^\sim = x^{\sim\sim} \rightsquigarrow y^-$;
- (psbck-c21) $(x \rightarrow y^{\sim\sim})^{\sim\sim} = x \rightarrow y^{\sim\sim}$ and $(x \rightsquigarrow y^{\sim\sim})^{\sim\sim} = x \rightsquigarrow y^{\sim\sim}$.

Proof

(psbck-c13) Since $0 \leq 0$, by $(psBCK_6)$ we get $0 \rightarrow 0 = 1$ and $0 \rightsquigarrow 0 = 1$, that is, $0^- = 1$ and $0^\sim = 1$.

Taking $x = 1$ and $y = 0$ in $(psBCK_2)$ we have $1 \leq (1 \rightarrow 0) \rightsquigarrow 0$, hence $(1 \rightarrow 0) \rightsquigarrow 0 = 1$. Thus by $(psBCK_6)$ we get $1 \rightarrow 0 \leq 0$, so $1 \rightarrow 0 = 0$, i.e. $1^- = 0$. Similarly, $1^\sim = 0$.

(psbck-c14) This follows by taking $y = 0$ in $(psBCK_2)$.

(psbck-c15) Applying $(psBCK_1)$ for $z = 0$ we get:

$$\begin{aligned} x \rightarrow y &\leq (y \rightarrow 0) \rightsquigarrow (x \rightarrow 0) = y^- \rightsquigarrow x^- \quad \text{and} \\ x \rightsquigarrow y &\leq (y \rightsquigarrow 0) \rightarrow (x \rightsquigarrow 0) = y^\sim \rightarrow x^\sim. \end{aligned}$$

(psbck-c16) From $x \leq y$, applying $(psbck-c1)$ we get $y \rightarrow 0 \leq x \rightarrow 0$, so $y^- \leq x^-$. Similarly, $y^\sim \leq x^\sim$.

(psbck-c17) By $(psbck-c15)$, $(psbck-c14)$ and $(psbck-c1)$ we get:

$$x \rightarrow y^\sim \leq y^{\sim\sim} \rightsquigarrow x^- \leq y \rightsquigarrow x^- \quad \text{and} \quad x \rightsquigarrow y^- \leq y^{\sim\sim} \rightarrow x^\sim \leq y \rightarrow x^\sim.$$

In the above inequalities we change x and y obtaining:

$$y \rightarrow x^\sim \leq x \rightsquigarrow y^- \quad \text{and} \quad y \rightsquigarrow x^- \leq x \rightarrow y^\sim.$$

Thus $x \rightarrow y^\sim = y \rightsquigarrow x^-$ and $x \rightsquigarrow y^- = y \rightarrow x^\sim$.

(*psbck-c18*) By (*psbck-c14*) and (*psbck-c16*) we get $x^{\sim\sim} \leq x^{\sim}$ and $x^{\sim\sim} \leq x^{-}$.

By (*psbck-c14*), replacing x with x^{\sim} and x^{-} we get $x^{\sim} \leq x^{\sim\sim}$ and $x^{-} \leq x^{\sim\sim}$, respectively. Thus $x^{\sim\sim} = x^{\sim}$ and $x^{\sim\sim} = x^{-}$.

(*psbck-c19*) By (*psbck-c17*) we have: $y \rightsquigarrow x^{-} = x \rightarrow y^{\sim}$.

Replacing y with y^{-} we get: $y^{-} \rightsquigarrow x^{-} = x \rightarrow y^{\sim\sim}$.

Replacing x by x^{\sim} in the last equality we get: $y^{-} \rightsquigarrow x^{\sim\sim} = x^{\sim} \rightarrow y^{\sim\sim}$.

Hence applying (*psbck-c18*) it follows that: $y^{-} \rightsquigarrow x^{-} = x^{\sim\sim} \rightarrow y^{\sim\sim}$.

Thus $x \rightarrow y^{\sim\sim} = y^{-} \rightsquigarrow x^{-} = x^{\sim\sim} \rightarrow y^{\sim\sim}$.

Similarly, $x \rightsquigarrow y^{\sim\sim} = y^{\sim} \rightarrow x^{\sim} = x^{\sim\sim} \rightsquigarrow y^{\sim\sim}$.

(*psbck-c20*) The assertions follow by replacing in (*psbck-c19*) y with y^{\sim} and y with y^{-} , respectively and applying (*psbck-c18*).

(*psbck-c21*) Applying (*psbck-c3*) and (*psbck-c19*) we have:

$$\begin{aligned} 1 &= (x \rightarrow y^{\sim\sim}) \rightsquigarrow (x \rightarrow y^{\sim\sim}) = x \rightarrow ((x \rightarrow y^{\sim\sim}) \rightsquigarrow y^{\sim\sim}) \\ &= x \rightarrow ((x \rightarrow y^{\sim\sim})^{\sim\sim} \rightsquigarrow y^{\sim\sim}) = (x \rightarrow y^{\sim\sim})^{\sim\sim} \rightsquigarrow (x \rightarrow y^{\sim\sim}). \end{aligned}$$

Hence $(x \rightarrow y^{\sim\sim})^{\sim\sim} \leq x \rightarrow y^{\sim\sim}$.

On the other hand, by (*psbck-c14*) we have $x \rightarrow y^{\sim\sim} \leq (x \rightarrow y^{\sim\sim})^{\sim\sim}$, thus $(x \rightarrow y^{\sim\sim})^{\sim\sim} = x \rightarrow y^{\sim\sim}$. Similarly, $(x \rightsquigarrow y^{\sim\sim})^{\sim\sim} = x \rightsquigarrow y^{\sim\sim}$. \square

We recall some notions and results regarding pseudo-BCK semilattices (see [209]).

Definition 1.3 A *pseudo-BCK join-semilattice* is an algebra $(A, \vee, \rightarrow, \rightsquigarrow, 1)$ such that (A, \vee) is a join-semilattice, $(A, \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK algebra and $x \rightarrow y = 1$ iff $x \vee y = y$.

Remark 1.2 It is easy to show that an algebra $(A, \vee, \rightarrow, \rightsquigarrow, 1)$ of type $(2, 2, 2, 0)$ is a pseudo-BCK join-semilattice if and only if (A, \vee) is a join-semilattice and $(A, \rightarrow, \rightsquigarrow, 1)$ satisfies (*psBCK'1*)–(*psBCK'5*) and the following identities:

$$(\text{psBCK}'_7) \quad x \vee [(x \rightarrow y) \rightsquigarrow y] = (x \rightarrow y) \rightsquigarrow y;$$

$$(\text{psBCK}'_8) \quad x \rightarrow (x \vee y) = 1.$$

Definition 1.4 A *pseudo-BCK meet-semilattice* is an algebra $(A, \wedge, \rightarrow, \rightsquigarrow, 1)$ such that (A, \wedge) is a meet-semilattice, $(A, \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK algebra and $x \rightarrow y = 1$ iff $x \wedge y = x$.

Remark 1.3 It is easy to show that an algebra $(A, \wedge, \rightarrow, \rightsquigarrow, 1)$ of type $(2, 2, 2, 0)$ is a pseudo-BCK meet-semilattice if and only if (A, \wedge) is a meet-semilattice and $(A, \rightarrow, \rightsquigarrow, 1)$ satisfies the identities (*psBCK'1*)–(*psBCK'5*) and the identities:

$$(psBCK_7'') \quad x \wedge [(x \rightarrow y) \rightsquigarrow y] = x;$$

$$(psBCK_8'') \quad (x \wedge y) \rightarrow y = 1.$$

Example 1.6 Given a pseudo-hoop $(A, \odot, \rightarrow, \rightsquigarrow, 1)$ (see Chap. 2), then $(A, \wedge, \rightarrow, \rightsquigarrow, 1)$ is a pseudo-BCK meet-semilattice, where $x \wedge y = x \odot (x \rightsquigarrow y) = (x \rightarrow y) \odot x$.

In the sequel by a *pseudo-BCK semilattice* we mean a pseudo-BCK join-semilattice.

Definition 1.5 Let $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ be a pseudo-BCK algebra. If the poset (A, \leq) is a lattice, then we say that A is a *pseudo-BCK lattice*.

A pseudo-BCK lattice is denoted by $(A, \wedge, \vee, \rightarrow, \rightsquigarrow, 1)$.

Example 1.7 Consider the bounded pseudo-BCK algebra $(A, \leq, \rightarrow, \rightsquigarrow, 0, 1)$ from Example 1.3. Since (A, \leq) is a lattice, it follows that A is a pseudo-BCK lattice.

Let A be a pseudo-BCK algebra. For all $x, y \in A$, define:

$$x \vee_1 y = (x \rightarrow y) \rightsquigarrow y, \quad x \vee_2 y = (x \rightsquigarrow y) \rightarrow y.$$

Proposition 1.5 *In any bounded pseudo-BCK algebra A the following hold for all $x, y \in A$:*

- (1) $0 \vee_1 x = x = 0 \vee_2 x$;
- (2) $x \vee_1 0 = x^{\rightsquigarrow\rightsquigarrow}, x \vee_2 0 = x^{\rightsquigarrow\rightsquigarrow}$;
- (3) $1 \vee_1 x = x \vee_1 1 = 1 = 1 \vee_2 x = x \vee_2 1$;
- (4) $x \leq y$ implies $x \vee_1 y = y$ and $x \vee_2 y = y$;
- (5) $x \vee_1 x = x \vee_2 x = x$.

Proof

- (1) $0 \vee_1 x = (0 \rightarrow x) \rightsquigarrow x = 1 \rightsquigarrow x = x$ and similarly $0 \vee_2 x = x$.
- (2) $x \vee_1 0 = (x \rightarrow 0) \rightsquigarrow 0 = x^{\rightsquigarrow\rightsquigarrow}$ and similarly $x \vee_2 0 = x^{\rightsquigarrow\rightsquigarrow}$.
- (3) We have: $1 \vee_1 x = (1 \rightarrow x) \rightsquigarrow x = 1$ and $x \vee_1 1 = (x \rightarrow 1) \rightsquigarrow 1 = 1$, so $1 \vee_1 x = x \vee_1 1 = 1$. Similarly, $1 \vee_2 x = x \vee_2 1 = 1$.
- (4) $x \vee_1 y = (x \rightarrow y) \rightsquigarrow y = 1 \rightsquigarrow y = y$. Similarly, $x \vee_2 y = y$.
- (5) This follows from the definitions of \vee_1 and \vee_2 . □

Proposition 1.6 *In any bounded pseudo-BCK algebra A the following hold for all $x, y \in A$:*

- (1) $x \vee_1 y^{\rightsquigarrow\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_1 y^{\rightsquigarrow\rightsquigarrow}$ and $x \vee_2 y^{\rightsquigarrow\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_2 y^{\rightsquigarrow\rightsquigarrow}$;
- (2) $x \vee_1 y^{\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_1 y^{\rightsquigarrow}$ and $x \vee_2 y^{\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_2 y^{\rightsquigarrow}$;
- (3) $(x^{\rightsquigarrow\rightsquigarrow} \vee_1 y^{\rightsquigarrow\rightsquigarrow})^{\rightsquigarrow\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_1 y^{\rightsquigarrow\rightsquigarrow}$ and $(x^{\rightsquigarrow\rightsquigarrow} \vee_2 y^{\rightsquigarrow\rightsquigarrow})^{\rightsquigarrow\rightsquigarrow} = x^{\rightsquigarrow\rightsquigarrow} \vee_2 y^{\rightsquigarrow\rightsquigarrow}$.

Proof

(1) Applying (*psbck-c19*) we have:

$$\begin{aligned} x \vee_1 y^{-\sim} &= (x \rightarrow y^{-\sim}) \rightsquigarrow y^{-\sim} = (x^{-\sim} \rightarrow y^{-\sim}) \rightsquigarrow y^{-\sim} = x^{-\sim} \vee_1 y^{-\sim}; \\ x \vee_2 y^{\sim-} &= (x \rightsquigarrow y^{\sim-}) \rightarrow y^{\sim-} = (x^{\sim-} \rightsquigarrow y^{\sim-}) \rightarrow y^{\sim-} = x^{\sim-} \vee_2 y^{\sim-}. \end{aligned}$$

(2) Applying (*psbck-c20*) we have:

$$\begin{aligned} x \vee_1 y^{\sim} &= (x \rightarrow y^{\sim}) \rightsquigarrow y^{\sim} = (x^{-\sim} \rightarrow y^{\sim}) \rightsquigarrow y^{\sim} = x^{-\sim} \vee_1 y^{\sim}; \\ x \vee_2 y^{-} &= (x \rightsquigarrow y^{-}) \rightarrow y^{-} = (x^{\sim-} \rightsquigarrow y^{-}) \rightarrow y^{-} = x^{\sim-} \vee_2 y^{-}. \end{aligned}$$

(3) Applying (*psbck-c21*) we have:

$$\begin{aligned} (x^{-\sim} \vee_1 y^{-\sim})^{-\sim} &= [(x^{-\sim} \rightarrow y^{-\sim}) \rightsquigarrow y^{-\sim}]^{-\sim} = (x^{-\sim} \rightarrow y^{-\sim}) \rightsquigarrow y^{-\sim} \\ &= x^{-\sim} \vee_1 y^{-\sim}; \\ (x^{\sim-} \vee_2 y^{\sim-})^{\sim-} &= [(x^{\sim-} \rightsquigarrow y^{\sim-}) \rightarrow y^{\sim-}]^{\sim-} = (x^{\sim-} \rightsquigarrow y^{\sim-}) \rightarrow y^{\sim-} \\ &= x^{\sim-} \vee_2 y^{\sim-}. \end{aligned} \quad \square$$

Proposition 1.7 *In any pseudo-BCK algebra the following hold for all $x, y \in A$:*

(*psbck-c22*) $(x \vee_1 y) \rightarrow y = x \rightarrow y$ and $(x \vee_2 y) \rightsquigarrow y = x \rightsquigarrow y$.

Proof This is a consequence of the property (*psbck-c11*). □

Lemma 1.1 *Let A be a pseudo-BCK algebra. Then:*

- (1) $x \vee_1 y$ ($y \vee_1 x$) is an upper bound of $\{x, y\}$;
- (2) $x \vee_2 y$ ($y \vee_2 x$) is an upper bound of $\{x, y\}$

for all $x, y \in A$.

Proof

(1) By (*psBCK₂*) we have $x \leq (x \rightarrow y) \rightsquigarrow y$.

Since by (*psbck-c6*), $y \leq (x \rightarrow y) \rightsquigarrow y$, we conclude that $x, y \leq x \vee_1 y$.

Similarly we get $x, y \leq y \vee_1 x$.

(2) Similar to (1). □

Definition 1.6 Let A be a pseudo-BCK algebra.

- (1) If $x \vee_1 y = y \vee_1 x$ for all $x, y \in A$, then A is called \vee_1 -commutative;
- (2) If $x \vee_2 y = y \vee_2 x$ for all $x, y \in A$, then A is called \vee_2 -commutative.

Lemma 1.2 *Let A be a pseudo-BCK algebra.*

- (1) If for all $x, y \in A$, $x \vee_1 y$ ($y \vee_1 x$) is the l.u.b. of $\{x, y\}$, then A is \vee_1 -commutative;
- (2) If for all $x, y \in A$, $x \vee_2 y$ ($y \vee_2 x$) is the l.u.b. of $\{x, y\}$, then A is \vee_2 -commutative.

Proof

- (1) Suppose that for all $x, y \in A$, $x \vee_1 y$ ($y \vee_1 x$) is the l.u.b. of $\{x, y\}$. Then by Lemma 1.1, for all $x, y \in A$ we have $y \vee_1 x \leq x \vee_1 y$ and $x \vee_1 y \leq y \vee_1 x$. Applying $(psBCK_5)$ we get $x \vee_1 y = y \vee_1 x$. Thus A is \vee_1 -commutative.
- (2) Similar to (1). □

Proposition 1.8 *Let A be a pseudo-BCK algebra.*

- (1) If A is \vee_1 -commutative, then $x \vee_1 y$ is the l.u.b. of $\{x, y\}$, for all $x, y \in A$;
- (2) If A is \vee_2 -commutative, then $x \vee_2 y$ is the l.u.b. of $\{x, y\}$, for all $x, y \in A$.

Proof

- (1) Let $x, y \in A$. According to Lemma 1.1, $x \vee_1 y$ is an upper bound of $\{x, y\}$. Let z be another upper bound of $\{x, y\}$, i.e. $x \leq z$ and $y \leq z$. We will prove that $x \vee_1 y \leq z$. Indeed, applying Proposition 1.5(4) and taking into consideration that A is \vee_1 -commutative we have:

$$\begin{aligned} x \vee_1 y \rightarrow z &= x \vee_1 y \rightarrow y \vee_1 z = x \vee_1 y \rightarrow z \vee_1 y \\ &= ((x \rightarrow y) \rightsquigarrow y) \rightarrow ((z \rightarrow y) \rightsquigarrow y). \end{aligned}$$

According to $(psBCK_1)$ we have $(b \rightarrow c) \rightsquigarrow (a \rightarrow c) \geq a \rightarrow b$ and replacing a with $z \rightarrow y$, b with $x \rightarrow y$ and c with y we get:

$$\begin{aligned} ((x \rightarrow y) \rightsquigarrow y) \rightarrow ((z \rightarrow y) \rightsquigarrow y) &\geq (z \rightarrow y) \rightsquigarrow (x \rightarrow y) \\ &\geq x \rightarrow z \quad (\text{by } (psBCK_1)). \end{aligned}$$

Hence $x \vee_1 y \rightarrow z \geq x \rightarrow z = 1$ (since $x \leq z$). It follows that $x \vee_1 y \rightarrow z = 1$, thus $x \vee_1 y \leq z$. We conclude that $x \vee_1 y$ is the l.u.b. of $\{x, y\}$.

- (2) Similar to (1). □

Theorem 1.1 *If A is a pseudo-BCK algebra, then:*

- (1) A is \vee_1 -commutative iff it is a join-semilattice with respect to \vee_1 (under \leq);
- (2) A is \vee_2 -commutative iff it is a join-semilattice with respect to \vee_2 (under \leq).

Proof This is a consequence of Lemma 1.2 and Proposition 1.8. □

Corollary 1.1 *Let A be a pseudo-BCK algebra. Then:*

- (1) If A is \vee_1 -commutative, then $x \vee_1 y \leq x \vee_2 y, y \vee_2 x$ for all $x, y \in A$;

(2) If A is \vee_2 -commutative, then $x \vee_2 y \leq x \vee_1 y, y \vee_1 x$ for all $x, y \in A$.

Proof

(1) According to Lemma 1.1, $x \vee_2 y, y \vee_2 x$ are upper bounds of $\{x, y\}$. By Proposition 1.8, $x \vee_1 y$ is the l.u.b. of $\{x, y\}$, thus $x \vee_1 y \leq x \vee_2 y, y \vee_2 x$.

(2) Similar to (1). \square

Definition 1.7 A pseudo-BCK algebra is called *sup-commutative* if it is both \vee_1 -commutative and \vee_2 -commutative.

Theorem 1.2 A pseudo-BCK algebra is *sup-commutative* iff it is a join-semilattice with respect to both \vee_1 and \vee_2 .

Proof This follows from Theorem 1.1. \square

Corollary 1.2 If A is a *sup-commutative pseudo-BCK algebra*, then $x \vee_1 y = x \vee_2 y$ for all $x, y \in A$.

Proof By Corollary 1.1, $x \vee_1 y \leq x \vee_2 y$ and $x \vee_2 y \leq x \vee_1 y$, hence $x \vee_1 y = x \vee_2 y$. \square

Lemma 1.3 In a \vee_1 -commutative (\vee_2 -commutative) bounded pseudo-BCK algebra A , we have $x^{\sim\sim} = x$ ($x^{\sim-} = x$, respectively), for all $x \in A$.

Proof Replacing y with 0 in the identity $x \vee_1 y = y \vee_1 x$, we get $(x \rightarrow 0) \rightsquigarrow 0 = (0 \rightarrow x) \rightsquigarrow x$, i.e. $x^{\sim\sim} = x$.

Similarly, replacing y with 0 in $x \vee_2 y = y \vee_2 x$, we get $x^{\sim-} = x$. \square

Corollary 1.3 Let A be a *sup-commutative, bounded pseudo-BCK algebra*. Then $x^{\sim\sim} = x^{\sim-} = x$, for all $x \in A$.

Proof This follows by replacing y with 0 in the equality $x \vee_1 y = x \vee_2 y$ and applying Lemma 1.3. \square

In a bounded pseudo-BCK algebra A , define, for all $x, y \in A$:

$$x \wedge_1 y := (x^- \vee_1 y^-)^{\sim},$$

$$x \wedge_2 y := (x^- \vee_2 y^-)^{\sim}.$$

Lemma 1.4 Let A be a pseudo-BCK algebra. Then for all $x, y \in A$:

(1) $x \wedge_1 y$ ($y \wedge_1 x$) is a lower bound of $\{x^{\sim\sim}, y^{\sim\sim}\}$;

(2) $x \wedge_2 y$ ($y \wedge_2 x$) is a lower bound of $\{x^{\sim-}, y^{\sim-}\}$.

Proof

- (1) By Lemma 1.1 we have $x^-, y^- \leq x^- \vee_1 y^-$, hence $x \wedge_1 y = (x^- \vee_1 y^-)^\sim \leq x^{-\sim}, y^{-\sim}$. Thus $x \wedge_1 y$ is a lower bound of $\{x^{-\sim}, y^{-\sim}\}$.
 (2) Similar to (1). □

Proposition 1.9 *Let A be a bounded pseudo-BCK algebra.*

- (1) *If A is \vee_1 -commutative, then $x \wedge_1 y$ ($y \wedge_1 x$) is the g.l.b. of $\{x, y\}$ and $x \wedge_1 y = y \wedge_1 x$, for all $x, y \in A$;*
 (2) *If A is \vee_2 -commutative, then $x \wedge_2 y$ ($y \wedge_2 x$) is the g.l.b. of $\{x, y\}$ and $x \wedge_2 y = y \wedge_2 x$, for all $x, y \in A$.*

Proof

- (1) By Lemma 1.3, $x^{-\sim} = x$ and $y^{-\sim} = y$. Hence by Lemma 1.4, $x \wedge_1 y$ is a lower bound of $\{x, y\}$. Now let z be another lower bound of $\{x, y\}$, i.e. $z \leq x, y$. It follows that $x^-, y^- \leq z^-$, thus z^- is an upper bound of $\{x^-, y^-\}$. Since A is \vee_1 -commutative, by Proposition 1.8, $x^- \vee_1 y^-$ is the l.u.b. of $\{x^-, y^-\}$, hence $x^- \vee_1 y^- \leq z^-$. Thus $z = z^{-\sim} \leq (x^- \vee_1 y^-)^\sim = x \wedge_1 y$, i.e. $x \wedge_1 y$ is the g.l.b. of $\{x, y\}$. Since A is \vee_1 -commutative, we have $x^- \vee_1 y^- = y^- \vee_1 x^-$, hence by definition it follows that $x \wedge_1 y = y \wedge_1 x$, for all $x, y \in A$.
 (2) Similar to (1). □

Corollary 1.4 *Let A be a bounded pseudo-BCK algebra.*

- (1) *If A is \vee_1 -commutative, then A is a lattice with respect to \wedge_1, \vee_1 ;*
 (2) *If A is \vee_2 -commutative, then A is a lattice with respect to \wedge_2, \vee_2 .*

Proof This follows by Propositions 1.8 and 1.9. □

Theorem 1.3 *A bounded sup-commutative pseudo-BCK algebra A is a lattice with respect to both \vee_1, \wedge_1 and \vee_2, \wedge_2 (under \leq) and for all x, y we have:*

$$x \vee_1 y = x \vee_2 y, \quad x \wedge_1 y = x \wedge_2 y.$$

Proof By Corollary 1.4, A is a lattice with respect to both \wedge_1, \vee_1 and \wedge_2, \vee_2 . By Corollary 1.2, $x \vee_1 y = x \vee_2 y$ for all $x, y \in A$. By Proposition 1.9 we get: $x \wedge_2 y \leq x \wedge_1 y$ and $x \wedge_1 y \leq x \wedge_2 y$, hence $x \wedge_1 y = x \wedge_2 y$ for all $x, y \in A$. □

We recall that a *downwards-directed set* (or a *filtered set*) is a partially ordered set (A, \leq) such that whenever $a, b \in A$, there exists an $x \in A$ such that $x \leq a$ and $x \leq b$.

Dually, an *upwards-directed set* is a partially ordered set (A, \leq) such that whenever $a, b \in A$, there exists an $x \in A$ such that $a \leq x$ and $b \leq x$.

If X is a set, then a *net* in X will be a set $\{x_i \mid i \in I\}$, where (I, \leq) is an upwards-directed set.

We say that a pseudo-BCK algebra A satisfies the *relative cancellation property*, (RCP) for short, if for every $a, b, c \in A$,

$$a, b \leq c \quad \text{and} \quad c \rightarrow a = c \rightarrow b, c \rightsquigarrow a = c \rightsquigarrow b \quad \text{imply} \quad a = b.$$

We note that a pseudo-BCK algebra A that is sup-commutative and satisfies the (RCP) condition is said to be a *Łukasiewicz pseudo-BCK algebra* (see [112]).

Example 1.8 The pseudo-BCK algebra A from Example 1.3 is downwards-directed with (RCP).

Proposition 1.10 *Any downwards-directed sup-commutative pseudo-BCK algebra has (RCP).*

Proof Consider $a, b, c \in A$ such that $a, b \leq c$ and $c \rightarrow a = c \rightarrow b, c \rightsquigarrow a = c \rightsquigarrow b$. There exists an $x \in A$ such that $x \leq a, b$.

By (*psbck-c₁*), from $a \leq c$ it follows that $c \rightsquigarrow x \leq a \rightsquigarrow x$.

According to Proposition 1.5(4) and (*psbck-c₃*) we have:

$$\begin{aligned} a \rightsquigarrow x &= (c \rightsquigarrow x) \vee_1 (a \rightsquigarrow x) = (a \rightsquigarrow x) \vee_1 (c \rightsquigarrow x) \\ &= [(a \rightsquigarrow x) \rightarrow (c \rightsquigarrow x)] \rightsquigarrow (c \rightsquigarrow x) = [c \rightsquigarrow [(a \rightsquigarrow x) \rightarrow x]] \rightsquigarrow (c \rightsquigarrow x) \\ &= [c \rightsquigarrow (a \vee_2 x)] \rightsquigarrow (c \rightsquigarrow x) = [c \rightsquigarrow (x \vee_2 a)] \rightsquigarrow (c \rightsquigarrow x) \\ &= (c \rightsquigarrow a) \rightsquigarrow (c \rightsquigarrow x). \end{aligned}$$

Similarly, $b \rightsquigarrow x = (c \rightsquigarrow b) \rightsquigarrow (c \rightsquigarrow x) = (c \rightsquigarrow a) \rightsquigarrow (c \rightsquigarrow x) = a \rightsquigarrow x$.

We have: $a = x \vee_2 a = a \vee_2 x = (a \rightsquigarrow x) \rightarrow x = (b \rightsquigarrow x) \rightarrow x = b \vee_2 x = x \vee_2 b = b$.

Thus A has (RCP). \square

1.2 Pseudo-BCK Algebras with Pseudo-product

Definition 1.8 A pseudo-BCK algebra with the (*pP*) condition (i.e. with the *pseudo-product* condition) or a *pseudo-BCK(pP) algebra* for short, is a pseudo-BCK algebra $(A, \leq, \rightarrow, \rightsquigarrow, 1)$ satisfying the (pP) condition:

(pP) For all $x, y \in A$, $x \odot y$ exists where

$$x \odot y = \min\{z \mid x \leq y \rightarrow z\} = \min\{z \mid y \leq x \rightsquigarrow z\}.$$

Example 1.9 Take $A = \{0, a_1, a_2, s, a, b, n, c, d, m, 1\}$ with $0 < a_1 < a_2 < s < a, b < n < c, d < m < 1$ (see Fig. 1.3).

Consider the operations $\rightarrow, \rightsquigarrow$ given by the following tables:

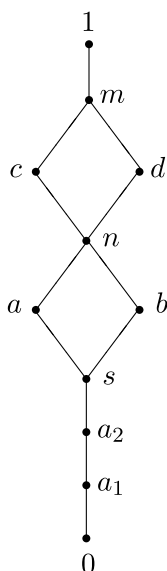


Fig. 1.3 Example of bounded pseudo-BCK(pP) algebra

\rightarrow	0	a_1	a_2	s	a	b	n	c	d	m	1
0	1	1	1	1	1	1	1	1	1	1	1
a_1	a_1	1	1	1	1	1	1	1	1	1	1
a_2	a_1	a_1	1	1	1	1	1	1	1	1	1
s	0	a_1	a_2	1	1	1	1	1	1	1	1
a	0	a_1	a_2	m	1	m	1	1	1	1	1
b	0	a_1	a_2	m	m	1	1	1	1	1	1
n	0	a_1	a_2	m	m	m	1	1	1	1	1
c	0	a_1	a_2	m	m	m	m	1	m	1	1
d	0	a_1	a_2	m	m	m	m	m	1	1	1
m	0	a_1	a_2	m	m	m	m	m	m	1	1
1	0	a_1	a_2	s	a	b	n	c	d	m	1
\rightsquigarrow	0	a_1	a_2	s	a	b	n	c	d	m	1
0	1	1	1	1	1	1	1	1	1	1	1
a_1	a_2	1	1	1	1	1	1	1	1	1	1
a_2	0	a_1	1	1	1	1	1	1	1	1	1
s	0	a_1	a_2	1	1	1	1	1	1	1	1
a	0	a_1	a_2	m	1	m	1	1	1	1	1
b	0	a_1	a_2	m	m	1	1	1	1	1	1
n	0	a_1	a_2	m	m	m	1	1	1	1	1
c	0	a_1	a_2	m	m	m	m	1	m	1	1
d	0	a_1	a_2	m	m	m	m	m	1	1	1
m	0	a_1	a_2	m	m	m	m	m	m	1	1
1	0	a_1	a_2	s	a	b	n	c	d	m	1